

Stake-governed random-turn games

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Tug of war

.... in ECONOMICS

Harris
&
Vickers
1987
'Racing
with
uncertainty...'
≈ 600 citations

.... and in MATHEMATICS

Peres, Schramm,
Sheffield & Wilson
2009:
'Tug of war and
the infinity Laplacian'
≈ 500 citations

Harmonic functions on a graph

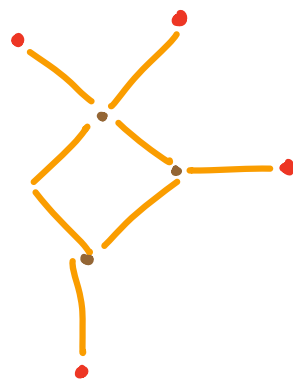
A graph $G = (V, E)$ has boundary $B \subseteq V$ with data $f: B \rightarrow \mathbb{R}$.

Suppose that $h: V \rightarrow \mathbb{R}$, $h|_B = f$, is such that, for each $v \in V \setminus B$,

$\omega = h(v) \in \mathbb{R}$ minimizes

$$\sum_{u \sim v} (\omega - h(u))^2.$$

Then h is harmonic — we may say ‘ 2 -harmonic’ —
and $h(v) = \frac{1}{\deg(v)} \sum_{u \sim v} h(u) \quad \forall v \in V \setminus B.$



2-harmonic functions and simple random walk

Simple random walk $X: \mathbb{N} \rightarrow V$, $x(0) = v \in V \setminus B$,
jumps at each step to a **uniformly chosen** neighbour.

Let $\tau =$ the time that X reaches B .

Then the 2-harmonic extension $h: V \rightarrow \mathbb{R}$, $h|_B = f$,
satisfies

$$h(v) = \mathbb{E} \left[f(x(\tau)) \mid x(0) = v \right].$$

We could call X the **2-walk** !

p-harmonic functions for $p \in (1, \infty)$

Say that $h: V \rightarrow \mathbb{R}$, $h|_B = f$, is *p-harmonic* if,

for each $v \in V \setminus B$, $h(v)$ equals
the minimizer w of

$$\left(\sum_{u \sim v} |w - h(u)|^p \right)^{1/p}.$$



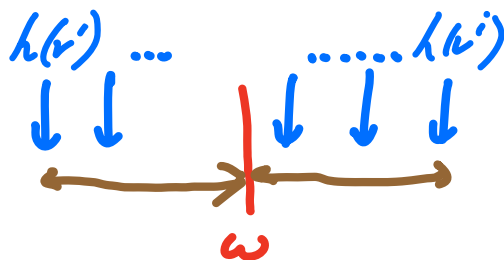
The case $p = \infty$.

Say that $h: V \rightarrow \mathbb{R}$, $h|_B = f$,
is *infinity harmonic* if,

for each $v \in V \setminus B$,

$h(v)$ equals the minimizer ω of

$$\max_{u \sim v} |\omega - h(u)|$$



Infinitely harmonic functions

In other words, h is ∞ -harmonic if

$$h(v) = \frac{1}{2} \left(\max_{u \sim v} h(u) + \min_{u \sim v} h(u) \right)$$

$$\forall v \in V \setminus B.$$

What then is the infinity walk $X: \mathbb{N} \rightarrow V$

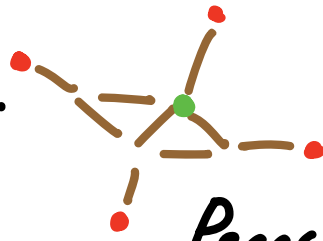
for which

$$h(v) = \mathbb{E} \left[f(X(\tau)) \mid X(0) = v \right]$$

(∞ -harmonic arrival of B)

The game RANDOM TUG-OF-WAR

A counter \bullet begins at vertex B .



Mina and **Maxine** win the first turn according to a fair coin flip.

Peres,
Schramm,
Sheffield,
Wilson 2007

The turn vector moves the counter to an adjacent vertex. Similarly for later turns.

The game ends when the counter reaches B , and MINA pays MAXINE

$$\text{Pay} = f(\text{terminal counter location}).$$

Strategies

A strategy for a given player specifies what move to play in any given game position.

Let \mathcal{S}_- = set of strategies for *Mina*
and \mathcal{S}_+ = set of strategies for *Maxine*

For $(s_-, s_+) \in \mathcal{S}_- \times \mathcal{S}_+$, set

$$M(s_-, s_+) = \mathbb{E}[\text{Pay}(s_-, s_+)]$$

to be the mean value of the terminal payment
Mina \rightarrow *Maxine* under the strategy pair (s_-, s_+) .

Game value

Suppose that **Mina** must declare to Maxine her strategy at the start of the game.

The BEST she can do is her game value

$$V_- = \inf_{s_- \in S_-} \sup_{s_+ \in S_+} u(s_-, s_+).$$

Similarly, if it is **Maxine** who must declare her strategy:

$$V_+ = \sup_{s_+ \in S_+} \inf_{s_- \in S_-} u(s_-, s_+).$$

$V_+ \leq V_-$ always. If $V_+ = V_-$, then the game
this is the game value has value.

Game value and optimal strategies in tug-of-war

Theorem [PSSW09]

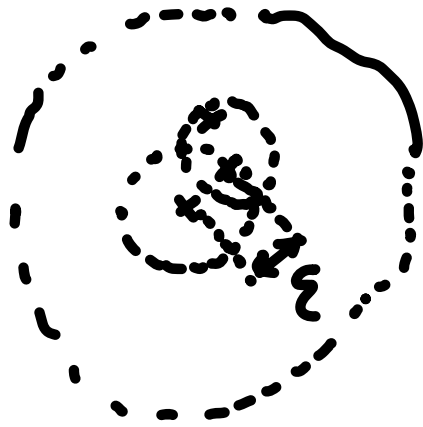
On any finite graph, the game value exists.
It is the infinity harmonic extension

$$h(v) = \frac{1}{2} \left(\max_{u \sim v} h(u) + \min_{u \sim v} h(u) \right)$$

of the boundary data $f: B \rightarrow \mathbb{R}$.

Under optimal play, **Max** plays to an h -maximizing neighbour; **Min** plays to an h -minimizing one.

Aside: continuum infinity harmonic functions



PSSW
→
2007

gone value
in low ε limit
is **CONTINUUM**
 ∞ -harmonic:

RANDOM TUG-OF-WAR
with ε -sized steps
in a domain in \mathbb{R}^d

$$0 = \Delta_{\infty} u := \sum_{i,j} u_{x_i} u_{x_i x_j} u_{x_j} .$$

Question: what is the STRATEGIC IMPORTANCE of a position in a multi-turn game?



Suppose that Maxine is offered the right to buy the first node, with the later game running as usual.

How much should she be prepared to pay?

If she buys the first move,



her mean payoff under jointly optimal

play is $\frac{n+1}{2n} = \frac{1}{2} + \frac{1}{2n}$.

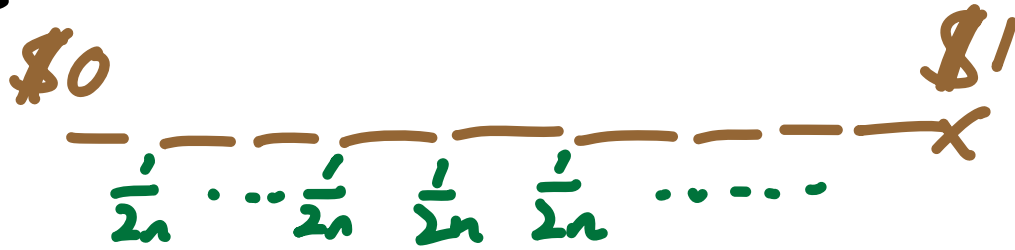
If she does not buy this move,

this payoff is $\frac{1}{2}$.

A horizontal sequence of nodes representing a game tree. The nodes are labeled from left to right as 0 , 1 , 2 , and $2n$. Above each node is an 'x' and below each node is a '0'. The node labeled 1 is circled in green, indicating it is the chosen move.


So the fair price of the purchase is $\frac{1}{2n}$.

Equally for any other starting position.



Apparent conclusion:

all positions are EQUALLY important,
and each is WORTH $\Theta\left(\frac{1}{n}\right)$.

Another view 
Under optimal play, the game lasts
 $\Theta(n^2)$ steps, because the counter X
follows SIMPLE RANDOM WALK.

If all STEPS are EQUALLY important,
and the total sum at stake is \$1,
then the first turn is WORTH $\Theta(\frac{1}{n^2})$.

Two differing predictions

x x x x x x x x x
0 n 2n

The importance of the first term is

$$\Theta\left(\frac{1}{n}\right)$$

$$\Theta\left(\frac{1}{n^2}\right)$$

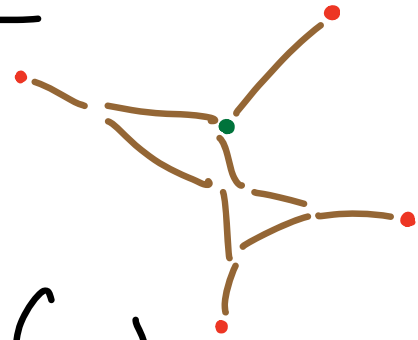
according to

the SPACE
prediction

the TIME
prediction

State-governed tug-of-war

Let $\lambda \in (0, \infty)$. Suppose that
Mina and *Maxine* have
respective FORTUNES of one and λ
at the OUTSET.



At the start of a given turn,
each retains some part of her original
fortune. Each is asked to stake some part
of this fortune — *Maxine* stakes a , *Mina* b .

These sums are deducted from the respective fortunes.

Maxine wins the RIGHT TO MOVE at the turn

with probability $\frac{a}{a+b}$; otherwise, Mia does.

The turn victor moves the counter to a neighbouring vertex, as in classical tug-of-war. The next turn proceeds, with the updated fortunes.

At the end of the game, Mia pays Maxine f (terminal counter location).

The players lose all of their remaining fortunes.

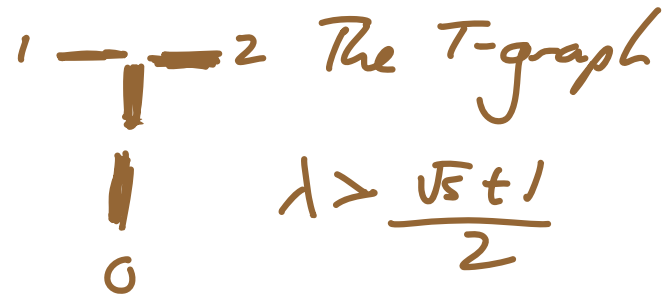
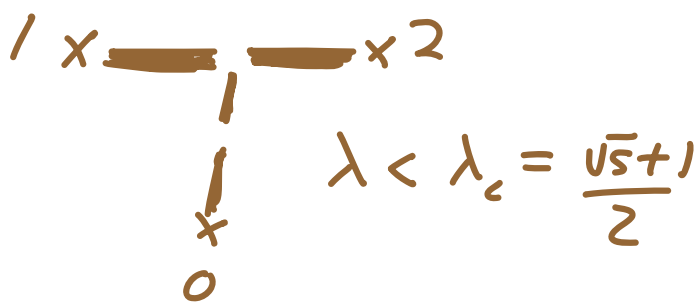
Two needed concepts

□ λ -biased ∞ -harmonic functions $h(\lambda, \cdot)$ satisfy

$$h(\lambda, v) = \frac{\lambda}{1+\lambda} \max_{u \sim v} h(\lambda, u) + \frac{1}{1+\lambda} \min_{u \sim v} h(\lambda, u),$$

$$h(\lambda, \cdot)|_B = f.$$

These may be computed by Peres-Šunić path decomposition



The T-graph

Two needed concepts

[2] Nash equilibrium

\mathcal{P}_- is the space of strategies for *Maria* ;
 \mathcal{P}_+ is this space for *Maxine*.

A pair $(s_-, s_+) \in \mathcal{P}_- \times \mathcal{P}_+$ is a Nash equilibrium if, for any $s'_- \in \mathcal{P}_-$ and $s'_+ \in \mathcal{P}_+$,

$$u(s_-, s'_+) \leq u(s_-, s_+) \leq u(s'_-, s_+).$$

At any Nash equilibrium, the payoff is the value of the game.

Dream-world game value

On a finite graph, set

$$\Delta(\lambda, v) = \max_{u \sim v} h(\lambda, u) - \min_{u \sim v} h(\lambda, u)$$

Proposition Suppose that $\Delta(\lambda, v) > 0 \forall v \in V \setminus B$.

Suppose that there is a pure Nash equilibrium in state-governed tug-of-war on the given graph.

Then the game value exists and equals the λ -biased α -harmonic function $h(\lambda, \cdot): V \rightarrow \mathbb{R}$ with $h(\lambda, \cdot)|_B = f$.

Proof (Stake strategy stealing)

Let (s_-^0, s_+^0) be an equilibrium.

Now let s_- be a strategy for Minia under which she stakes the same proportion (of her remaining fortune) as Maxine does under s_+^0 , and proposes to move to an h -minimizing neighbour.

Then the resulting $h(x, x_i)$ is a bounded supermartingale whose limiting value is the terminal payoff, because the game almost surely ends in finite time.

Thus, $\mu(s_-, s_+^o) \leq h(x_0) = h(v)$

But: (s_-, s_+^o) is an equilibrium \implies

$$\mu(s_-, s_+^o) \leq \mu(s_-, s_+^o).$$

Considering a similar strategy for *Maxine*
leads to

$$\mu(s_-, s_+^o) \geq h(v).$$

□

This conditional result prompts

TWO QUESTIONS :

1. Does a pure Nash equilibrium exist ?
2. If it does, then, when Maxine's fortune is t times Mia's, and the counter is at $v \in V \setminus B$, the players stake a shared proportion of their present fortunes — call this $S(t, v)$.

What then is the STAKE FUNCTION

$$V \setminus B \rightarrow \mathbb{R} : v \rightarrow S(t, v) ?$$

A plausible argument that identifies $S(\lambda, v)$

Mina has fortune 1, Maxine $\lambda \in (0, \infty)$.

The counter is at $v \in VLB$.

Write $S = S(\lambda, v)$.

Under optimal play, Maxine stakes λS
and Mina S .

What if Maxine deviates slightly,
to offer $(\lambda + \gamma)S$, with $0 < \gamma \ll 1$
at the first turn, with optimal later play
?

Change in near terminal payment

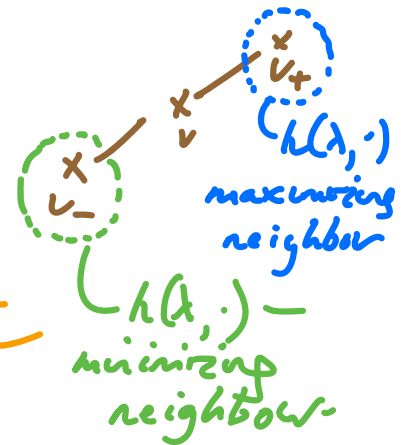
$$= \underbrace{(G - L)}_{\substack{\text{gain} \\ \text{(for Maxine)}}} \gamma + O(\gamma^2)$$

(loss)

The gain term equals

$$\zeta = \underbrace{\gamma^{-1} \lambda}_{\text{reflects the gain in probability of MAXINE winning the first turn}} \times \underbrace{\Delta(A, v)}_{= h(A, v_+) - h(A, v_-)}$$

reflects the gain in probability of MAXINE winning the first turn



It's tougher for *Maxine* at turns two and later:
her relative fortune is

$$\lambda_{alt} = \frac{\lambda - \lambda(s+1)}{1-s} = \lambda - \frac{\lambda s}{1-s} + O(s^2).$$

Thus, the loss term L takes the form

$$L = (1-s)^{-1} \frac{1}{(1+\lambda)^2} \mathbb{E} \sum_{i=1}^{\tau-1} \Delta(\lambda, X_i)$$

= game finish

Under optimal play,

$$Q = L.$$

This leads to a formula for the

STAKE function $S = S(\lambda, \nu)$:

$$S = \frac{\Delta(\lambda, \nu)}{\tau - 1}$$

← the space prediction

$$\mathbb{E} \sum_{i=0}^{\tau-1} \Delta(\lambda, x_i)$$

← similar to the time prediction

Alternatively, we may similarly predict:

$$J = \frac{\Delta(\lambda, \nu)}{(\lambda+1)^2 \frac{\partial}{\partial \lambda} h(\lambda, \nu)}$$

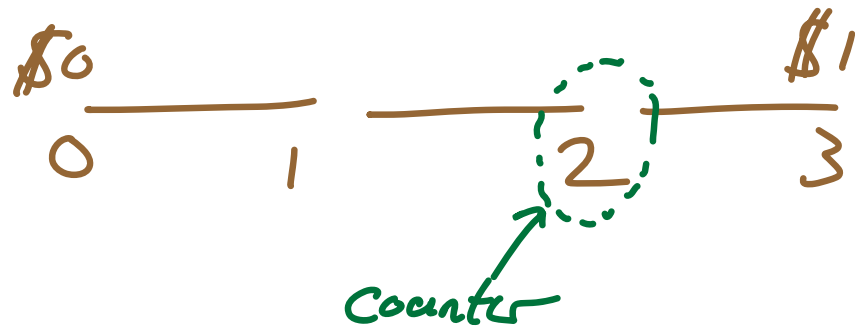
This is *space* versus *money*, or
short-term versus *long-term*.

A plausible picture, but ...

TWO PROBLEMS

1. For some SIMPLE graphs,
there is NO pure Nash equilibrium.
2. For some others, $h(G, v)$
fails to be differentiable in λ .

PROBLEM 1 : no pure Nash equilibrium

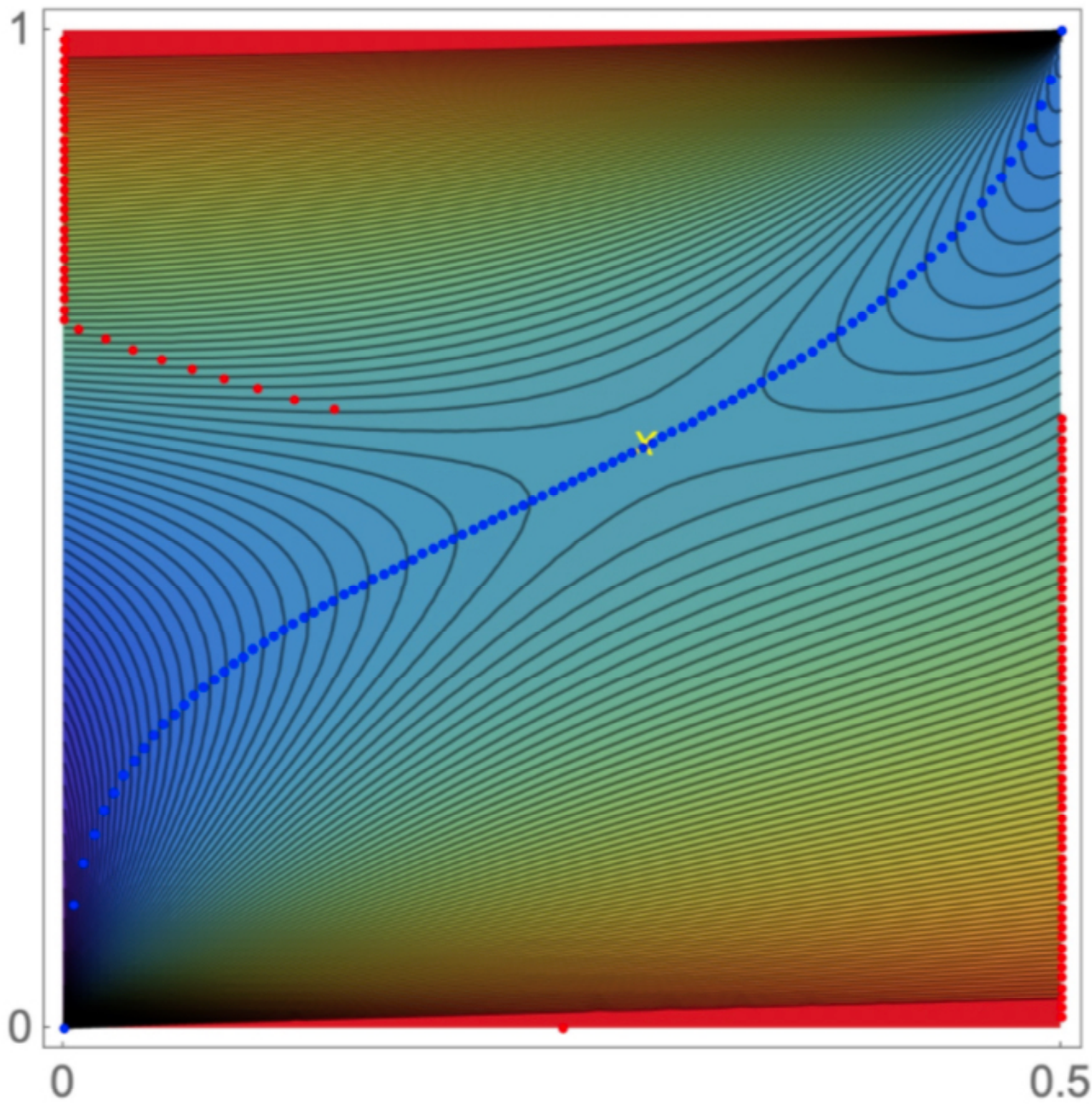


Maxine's relative fortune = $\frac{1}{2}$.

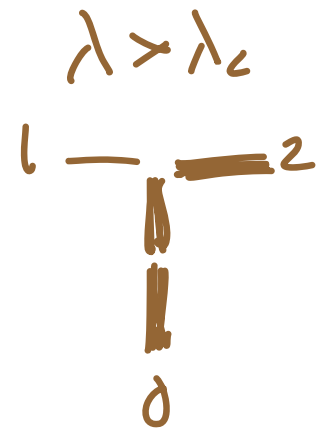
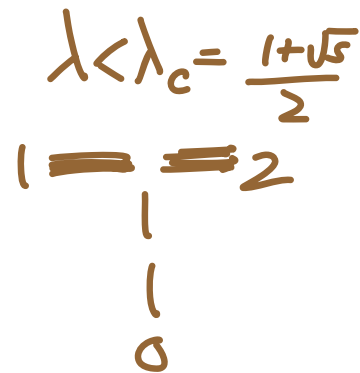
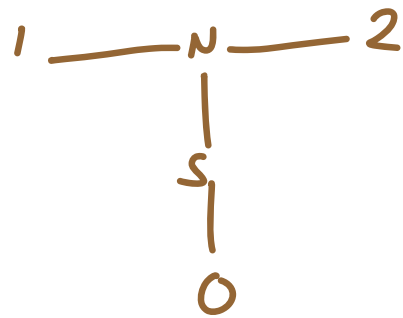
If a pure Nash equilibrium exists,
we can plot game value in the case that

Maxine stakes $a \in [0, \frac{1}{2}]$, and

Mia stakes $b \in [0, 1]$.



PROBLEM 2: $\lambda \rightarrow h(\lambda, \nu)$ may not be differentiable:



As λ rises through λ_c , the
 Peres-Šunić path decomposition changes,
 so that $\lambda \rightarrow h(\lambda, \nu)$ fails to be
 differentiable at λ_c .

In the proposed STAKE formula,

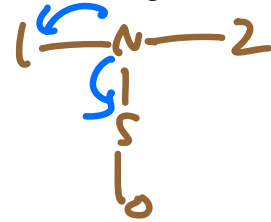
$$J = \frac{\Delta(A, v)}{\text{Denom}},$$

Denom equals $(\lambda+1)^2 \frac{\partial}{\partial \lambda} h(A, v)$ or $(\lambda+1)^2 \sum_{i=0}^{T-1} \Delta(x_i; \lambda)$.

But Denom is sometimes badly defined:

— h is not regular enough; or

— Minia can force different gameplays X
by proposing different moves



In summary, our TWO PROBLEMS may be construed as:

1. One or other player may be tempted to 'go for broke', so that a local saddle point fails to be global.
2. The λ -biased co-harmonic PATH decomposition may switch

as λ rises:

| | | | | | |
|---|---|---|---|---|-----------------------|
| 1 | — | N | — | 2 | |
| | | | | | |
| | | S | | | $\lambda < \lambda_c$ |
| | | | | | |
| | | 0 | | | |

| | | | | | |
|---|---|---|---|---|-----------------------|
| 1 | — | N | — | 2 | |
| | | | | | |
| | | S | | | $\lambda > \lambda_c$ |
| | | | | | |
| | | 0 | | | |

TWO MODIFICATIONS

that we now make in an effort to address the respective difficulties.

1. The game will become LEISURELY.
2. We will consider only a class of graphs called
ROOT-REWARD TREES.

Modification 1: the biased game.

We introduce a parameter $\varepsilon \in (0, 1]$

(to join λ , the relative fortune,
and v , the counter location).

RULE CHANGE:

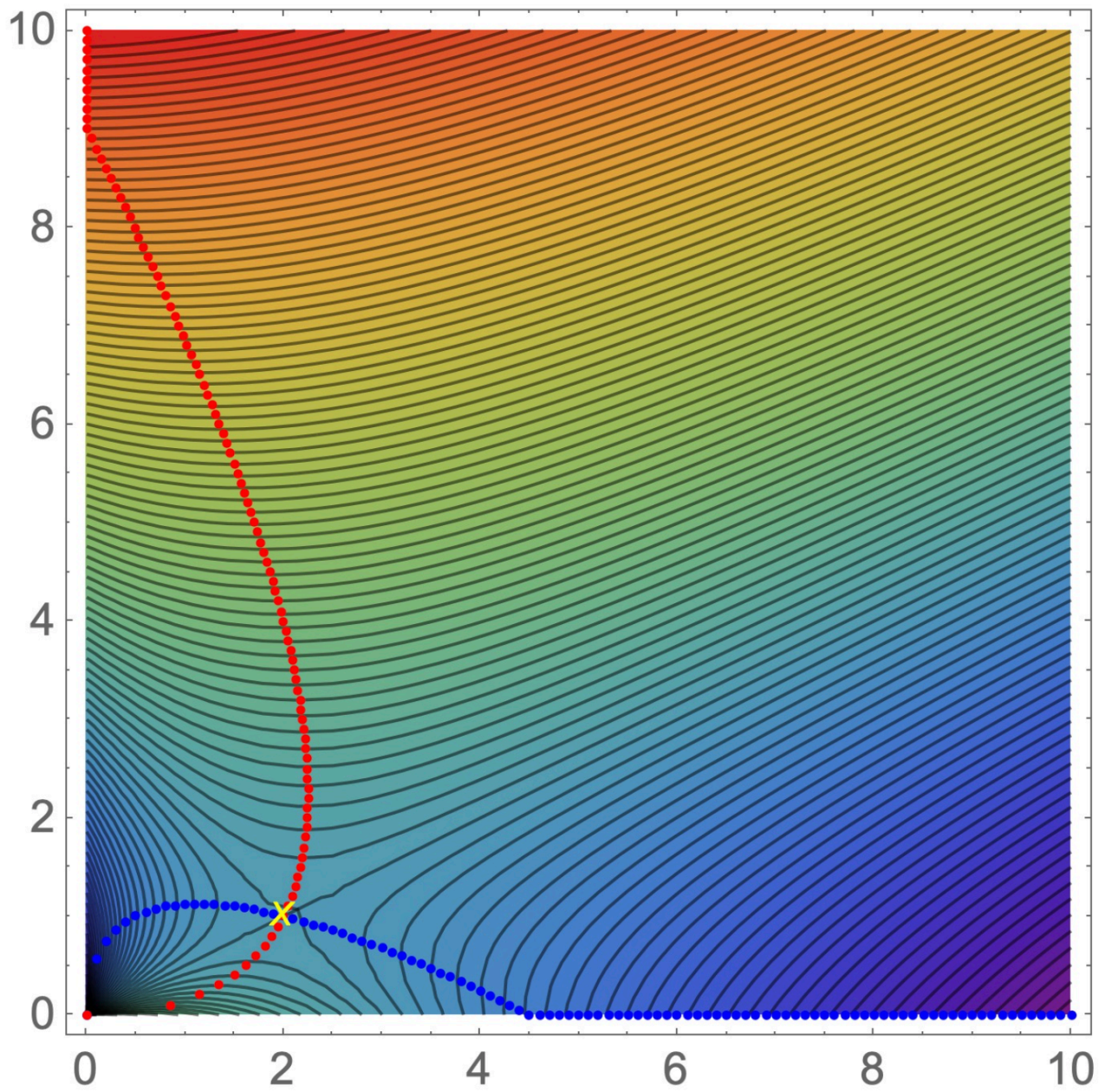
after players stake at a given turn,
the casino flips a coin that
lands HEADS with probability ε .

If the outcome is HEADS, then the turn proceeds as it would have done originally.

If it is TAILS, then no move takes place — the submitted stakes are simply lost to the players, and the counter moves nowhere.

IDEA: low choices of ϵ
disable the go-for-broke strategy.

would you bet your life-savings
when, more likely than not, they will
simply be swept from you ?



Modification 2: root-reward trees

A root-reward tree is

a finite tree $G = (V, E)$,

with boundary $B =$ set of leaves,

and $f: B \rightarrow \{0, 1\}$, $f = \mathbb{1}_r$

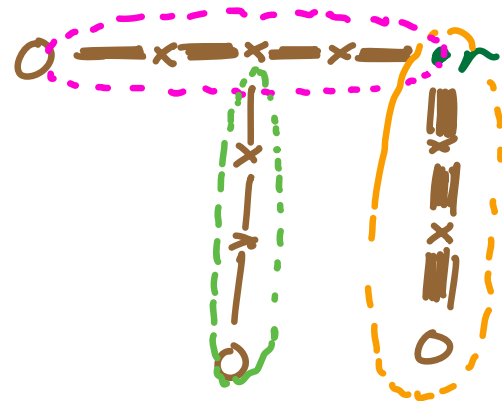
where r is a distinguished leaf called the root.

KEY FACT: the Peres-Sunic path decomposition of a λ -biased ω -harmonic function on a ROOT-REWARD tree is INDEPENDENT of $\lambda \in (0, \infty)$.

Root-reward tree and its 'essence' tree.



compute this
first
then this
and then this



Stage 1
2
3

This decomposition is determined by subtree diameters and is independent of $\lambda \in (0, \infty)$.

Game value and Nash equilibria for the leisurely stake game on root-reward trees

Theorem [H, Pete]

On any root-reward tree, and for any compact $K \subset (0, \infty)$, there is $\varepsilon_K > 0$ such that, for any relative fortune $\lambda \in K$, and for $0 < \varepsilon < \varepsilon_K$, in the leisurely game with parameter ε starting at any vertex $v \in V \setminus B$,

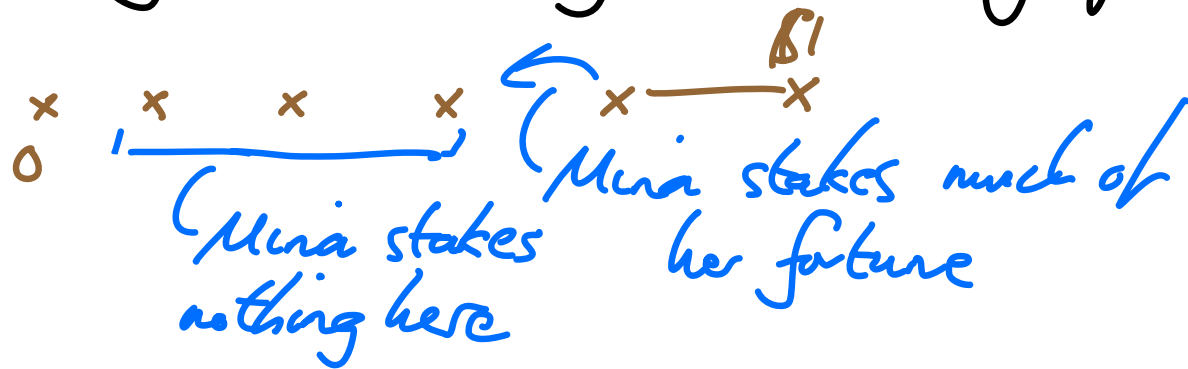
game value equals $h(\lambda, v)$;
every Nash equilibrium consists of
 $h(\lambda, \cdot)$ -maximizing/minimizing moves;
and the stakes are $(S, \lambda S)$, with

$$S = \frac{\varepsilon \Delta(\lambda, v)}{(\lambda+1) \frac{d}{d\lambda} h(\lambda, v)} = \frac{\Delta(\lambda, v)}{\mathbb{E} \sum_{i=0}^{\infty} \Delta(\lambda, x_i)}, \quad x_0 = v.$$

Caveat:

The payment *Mina* \rightarrow *Maxine* when play NEVER ENDS must be the FULL AMOUNT \$1 in this theorem.

Otherwise, *Mina* can fight *Maxine's* supposedly optimal play on a line graph



By so doing, *Mina* can force
infinite play.

So the FULL PAYMENT

Mina $\xrightarrow{\$1}$ *Maxine*

for infinite games is needed
to discourage *Mina* from doing this.

Self-funded state-governed tug-of-war.

In the game that we have considered, FORTUNES were limited, making them a PRECIOUS resource that players must spend — and conserve — over the likely lifetime of the game.

What about another means of making
a resource PRECIOUS ?

Both PLAYERS can spend whatever
they please at any given turn, but
now it is THEIR OWN MONEY that they spend
— so that the stakes of a given player
sum up to act as a running cost, to
be deducted from the terminal payment.

Self-funded stake-gained tug-of-war,
on finite line graphs and on \mathbb{Z} ,
is the subject of

'On the Trail of Lost Pennies',

arXiv: 2209.07451

— non-unique Nash equilibria on L_6



— countably many equilibria on \mathbb{Z}



