

# Asymptotic properties of weighted recursive and preferential attachment trees

ArXiv:1904.07115

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Retreat for Young Researchers in Probability and areas of Application,  
September 28, 2019

University of British Columbia

# Introduction

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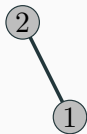
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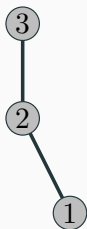
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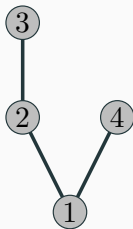
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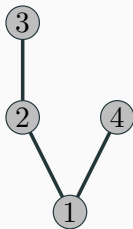


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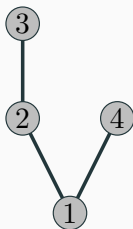
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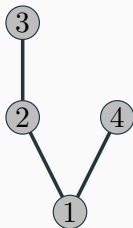
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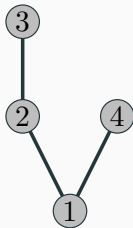


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- We can also use *random* sequences of weights.

**Goal:** study the asymptotic behaviour of some statistics of the tree  $T_n$  when  $n$  is large such as degrees, height of the tree, profile (number of vertices at any given height)...

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- Also, connection to the monkey walk (Mailler-Urbe 2018+).

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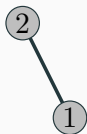
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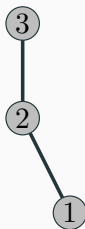
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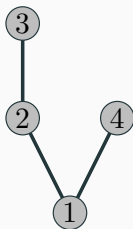
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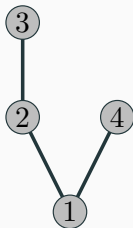
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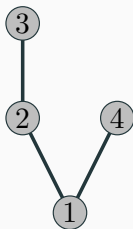


$$\forall k \in \{1, \dots, n\}, \quad \mathbb{P} \left( (n+1) \rightarrow k \mid P_n \right) = \frac{a_k + \text{deg}_n^+(k)}{\sum_{i=1}^n a_i + \sum_{i=1}^n \text{deg}_n^+(i)}.$$



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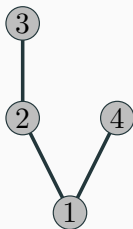
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## Theorem (S. 2019+)

*Preferential attachment trees are weighted recursive trees i.e. for any sequence of fitnesses  $\mathbf{a} = (a_n)_{n \geq 1}$ , there exists a random sequence  $(w_n^{\mathbf{a}})_{n \geq 1}$  such that the distributions  $\text{PA}((a_n)_{n \geq 1})$  and  $\text{WRT}((w_n^{\mathbf{a}})_{n \geq 1})$  coincide.*

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- If  $A_n \simeq c \cdot n$  as  $n \rightarrow \infty$  then the sequence  $(w_n^{\mathbf{a}})_{n \geq 1}$  satisfies  $W_n^{\mathbf{a}} \simeq \text{cst} \cdot n^\gamma$  a.s. with  $\gamma = \frac{c}{c+1}$ .

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- Results for weighted recursive trees automatically apply to preferential attachment trees!

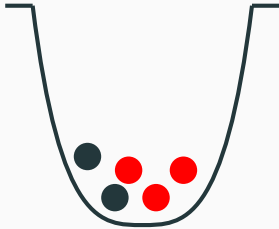
## DEFINITION OF THE RANDOM SEQUENCE $(w_n^a)_{n \geq 1}$

For any sequence  $\mathbf{a} = (a_n)_{n \geq 1}$  of fitnesses, the random sequence  $(w_n^a)_{n \geq 1}$  is defined as

$$w_1^a = W_1^a = 1 \quad \text{and} \quad \forall n \geq 2, \quad W_n^a = \prod_{k=1}^{n-1} \beta_k^{-1},$$

where the  $(\beta_k)_{k \geq 1}$  are independent with respective distribution  $\text{Beta}(A_k + k, a_{k+1})$ , with  $A_k = \sum_{i=1}^k a_i$ .

# THE CLASSICAL PÓLYA URN



- At time 0, the urn contains  $a$  black balls and  $b$  red balls.
- At each time  $n \geq 1$ , a ball is drawn from the urn and returned to the urn together with a new ball of the same colour.
- For all  $n \geq 1$ , we let  $X_n := \mathbf{1}_{\{\text{the ball drawn at time } n \text{ is black}\}}$ .

## Theorem

*Almost surely*

$$\frac{\#\{\text{black balls at time } n\}}{\#\{\text{balls at time } n\}} = \frac{a + \sum_{i=1}^n X_i}{a + b + n} \xrightarrow{n \rightarrow \infty} \beta,$$

where  $\beta$  is a random variable with distribution  $\text{Beta}(a, b)$ .



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where  $\beta$  is a random variable with distribution  $\text{Beta}(a, b)$ .

Furthermore, conditionally on  $\beta$ , the sequence  $(X_n)_{n \geq 1}$  is a sequence of i.i.d. Bernoulli random variables with parameter  $\beta$ .

## Convergence results for weighted recursive trees

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Proposition (S. 2019+)

If  $W_n \sim C \cdot n^\gamma$  as  $n \rightarrow \infty$  for  $0 < \gamma < 1$ , then we almost surely have

$$n^{-(1-\gamma)} \cdot (\deg_n^+(\textcircled{1}), \deg_n^+(\textcircled{2}), \dots) \xrightarrow[n \rightarrow \infty]{} \frac{1}{C(1-\gamma)} \cdot (w_1, w_2, \dots),$$

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- Thanks to the representation theorem, this also holds for preferential attachment trees.

For any  $k \geq 1$  we can write:

$$\text{deg}_n^+(\textcircled{k}) = \sum_{i=k+1}^n 1 \left\{ \textcircled{i} \rightarrow \textcircled{k} \right\}.$$

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The events  $\{\textcircled{i} \rightarrow \textcircled{k}\}$  for  $i \in \{k+1, k+2, \dots, n\}$  are *independent* and have respective probability  $\frac{w_k}{w_{i-1}}$ .

## CONVERGENCE OF DEGREES: ELEMENTS OF PROOFS

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The events  $\left\{ \textcircled{i} \rightarrow \textcircled{k} \right\}$  for  $i \in \{k+1, k+2, \dots, n\}$  are *independent* and have respective probability  $\frac{W_k}{W_{i-1}}$ . Hence by a law of large numbers, with probability 1 we have

$$\begin{aligned} \text{deg}_n^+(\textcircled{k}) &\underset{n \rightarrow \infty}{\sim} \sum_{i=k+1}^n \mathbb{P} \left( \textcircled{i} \rightarrow \textcircled{k} \right) \underset{n \rightarrow \infty}{\sim} \sum_{i=k+1}^n \frac{W_k}{W_{i-1}} \underset{n \rightarrow \infty}{\sim} W_k \cdot \sum_{i=k+1}^n \frac{1}{C \cdot i^\gamma} \\ &\underset{n \rightarrow \infty}{\sim} W_k \cdot \frac{n^{1-\gamma}}{C(1-\gamma)}. \end{aligned}$$



If  $(P_n)_{n \geq 1}$  is a sequence of trees with distribution  $\text{PA}((a_n)_{n \geq 1})$  with  $A_n \simeq c \cdot n$  as  $n \rightarrow \infty$  then we have the almost sure convergence in some  $\ell^p$  space,

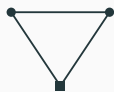
$$n^{-\frac{1}{c+1}} \cdot (\text{deg}_n^+(\textcircled{1}), \text{deg}_n^+(\textcircled{2}), \dots) \xrightarrow{n \rightarrow \infty} (m_1, m_2, \dots),$$

where  $(m_n)_{n \geq 1}$  is a constant times the sequence  $(w_n^a)_{n \geq 1}$ .

## Scaling limits for generalisation of Rémy's algorithm

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# A GENERALISED VERSION OF RÉMY'S ALGORITHM



$G_1$



$G_2$



$G_3$

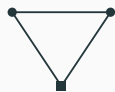


$G_4$



$G_5$

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$G_2$



$G_3$



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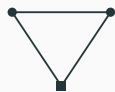


$G_5$



$H_1$

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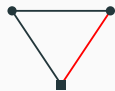
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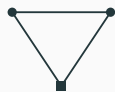
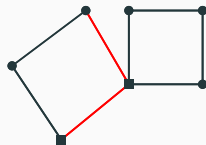
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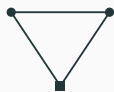
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 $G_1$  $G_2$  $G_3$  $G_4$  $G_5$ 

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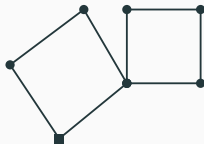
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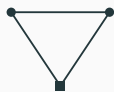


$G_5$



$H_2$

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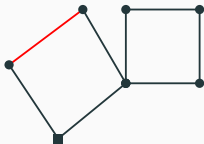
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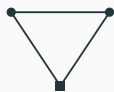


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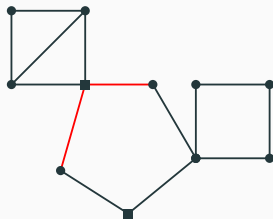
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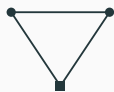
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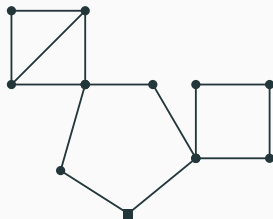
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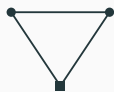
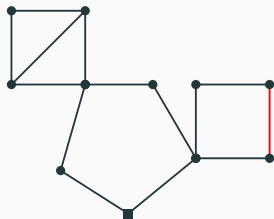


$G_5$

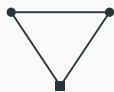
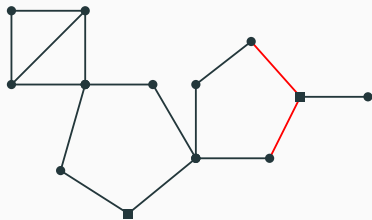


$H_3$

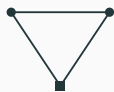
# A GENERALISED VERSION OF RÉMY'S ALGORITHM

 $G_1$  $G_2$  $G_3$  $G_4$  $G_5$ 

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 $G_1$  $G_2$  $G_3$  $G_4$  $G_5$ 

# A GENERALISED VERSION OF RÉMY'S ALGORITHM



$G_1$



$G_2$



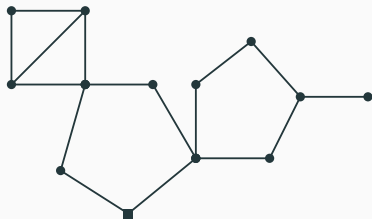
$G_3$



$G_4$

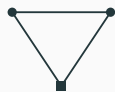


$G_5$



$H_4$

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$G_1$



$G_2$



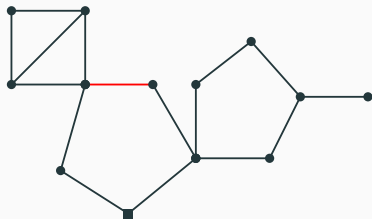
$G_3$



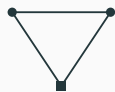
$G_4$



$G_5$



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$G_1$



$G_2$



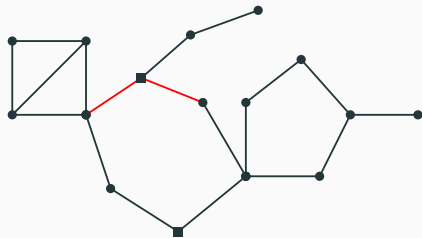
$G_3$



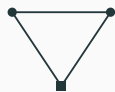
$G_4$



$G_5$



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$G_1$



$G_2$



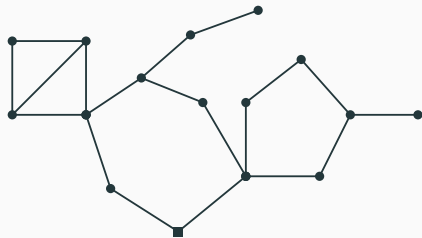
$G_3$



$G_4$



$G_5$



$H_5$



- In the original algorithm, the diameter of  $H_n$  grows as  $n^{1/2}$ , and we have an almost sure convergence in a metric space sense

$$(H_n, n^{-1/2} \cdot d_{gr}) \xrightarrow{n \rightarrow \infty} \mathcal{T},$$

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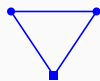
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- Do we get a scaling limit?
- If yes, what does it look like?

# A GENERALISED VERSION OF RÉMY'S ALGORITHM



$G_1$



$G_2$



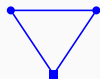
$G_3$



$G_4$



$G_5$



$H_1$



$T_1$

# A GENERALISED VERSION OF RÉMY'S ALGORITHM



$G_1$



$G_2$



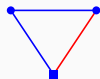
$G_3$



$G_4$



$G_5$



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$G_1$



$G_2$



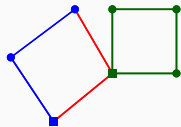
$G_3$



$G_4$



$G_5$



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$G_1$



$G_2$



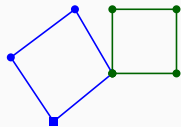
$G_3$



$G_4$



$G_5$



$H_2$



$T_2$



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$G_1$



$G_2$



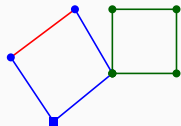
$G_3$



$G_4$



$G_5$



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$G_1$



$G_2$



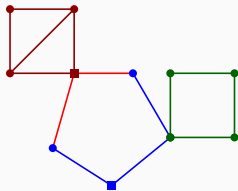
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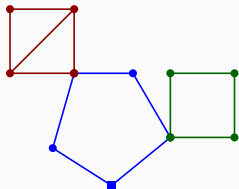
$G_3$



$G_4$



$G_5$



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$T_3$

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$G_1$



$G_2$



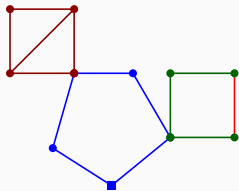
$G_3$



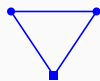
$G_4$



$G_5$



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$G_1$



$G_2$



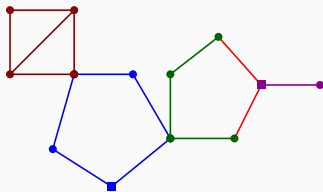
$G_3$



$G_4$



$G_5$



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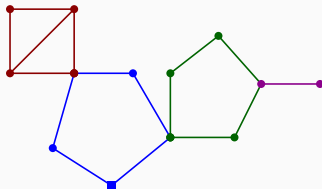
$G_3$



$G_4$



$G_5$



$H_4$



$T_4$

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$G_1$



$G_2$



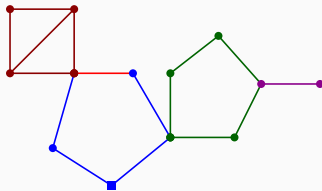
$G_3$



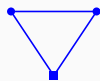
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$G_2$



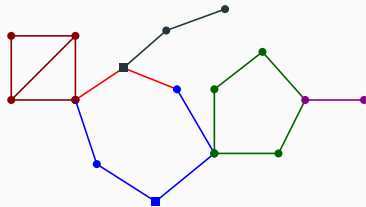
$G_3$



$G_4$



$G_5$





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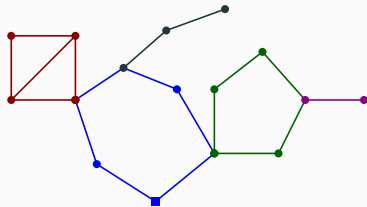
$G_3$



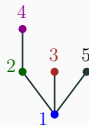
$G_4$



$G_5$



$H_5$



$T_5$

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$G_1$



$G_2$



$G_3$

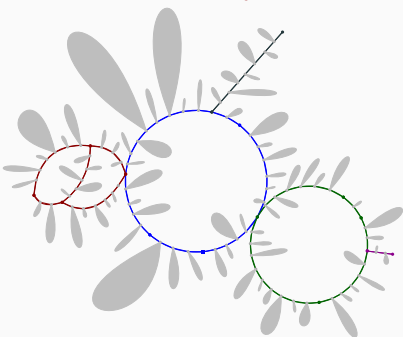


$G_4$



$G_5$

large  $n$   
distances  $\times n^{-1/(c+1)}$



## SCALING LIMIT OF THIS PROCESS

Denote by  $(a_n)_{n \geq 1} = (|E(G_n)|)_{n \geq 1}$  the sequence corresponding to the number of edges in  $(G_n)_{n \geq 1}$ .

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**Theorem (S. 2019+)**

Suppose that  $\sum_{i=1}^n a_i = c \cdot n + O(n^{1-\epsilon})$  for some constant  $c > 0$  and  $a_n \leq n^{\frac{1}{c+1}-\epsilon+o(1)}$ . Then we have the following convergence

$$\left( H_n, n^{-\frac{1}{c+1}} \cdot d_{\text{gr}}, \mu_{n, \text{unif}} \right) \xrightarrow[n \rightarrow \infty]{} (\mathcal{H}, d, \mu),$$

*almost surely in Gromov-Hausdorff-Prokhorov topology.*

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- Convergence of degrees in the tree  $\rightarrow$  convergence of each coloured portion of the graph.
- Controlling the degrees in the tree  $\rightarrow$  controlling distances in  $H_n$ .
- Description of the tree as a WRT  $\rightarrow$  iterative gluing construction.

# ITERATIVE GLUING CONSTRUCTION



$G_1$



$G_2$



$G_3$



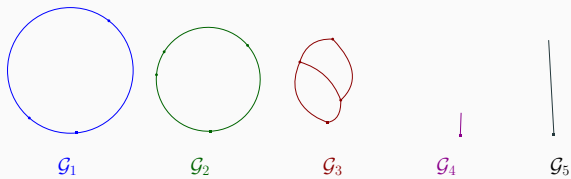
$G_4$



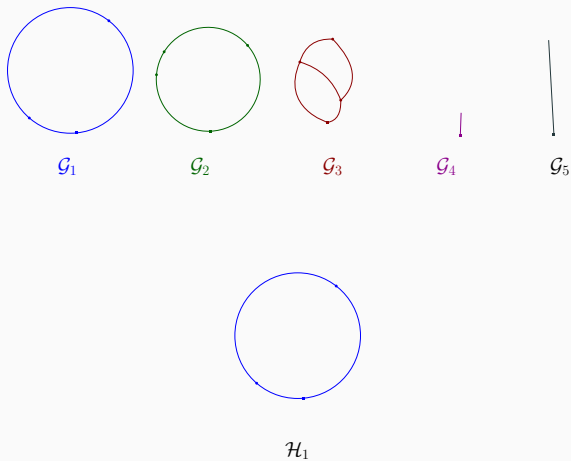
$G_5$



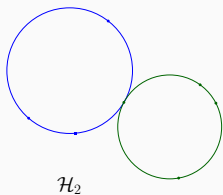
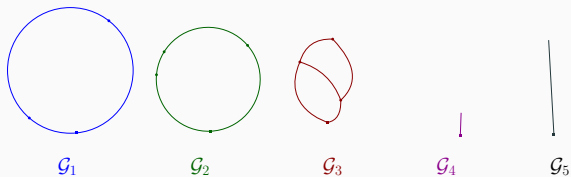
# ITERATIVE GLUING CONSTRUCTION



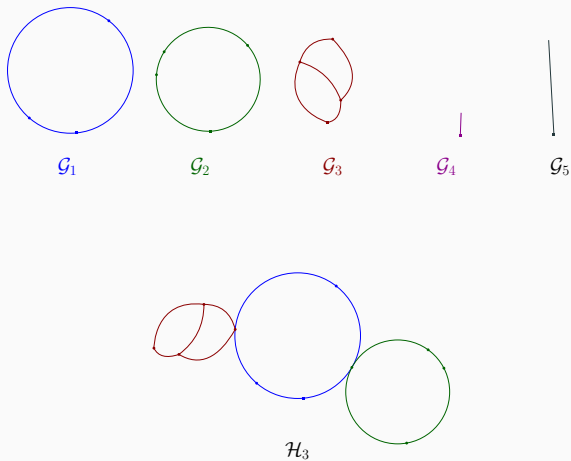
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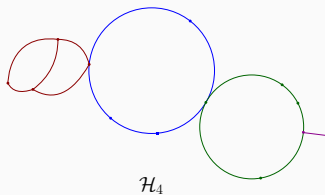
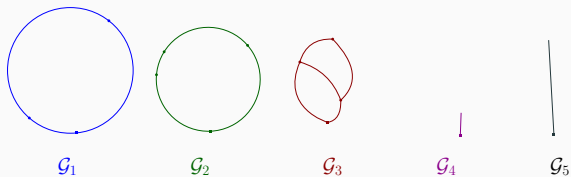
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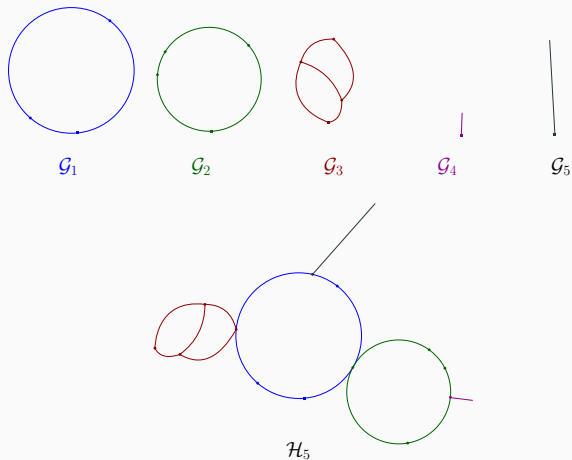
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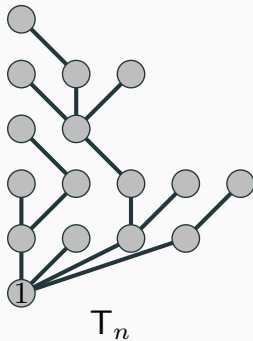


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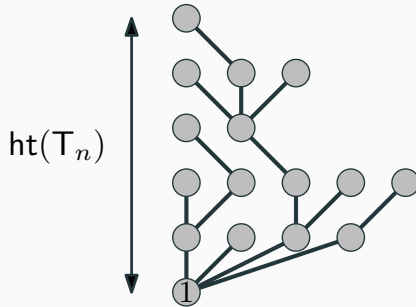
Thank you for your attention!

# HEIGHT AND PROFILE OF WRT

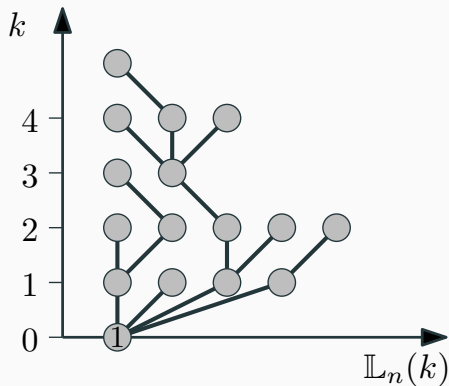




# HEIGHT AND PROFILE OF WRT



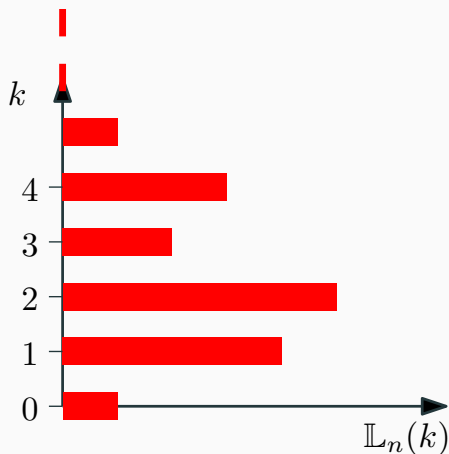
## HEIGHT AND PROFILE OF WRT



The profile of the tree  $T_n$  is the function  $\mathbb{L}_n : \mathbb{N} \rightarrow \mathbb{N}$  defined as

$$\forall k \geq 0, \quad \mathbb{L}_n(k) = \#\{\text{vertices at height } k \text{ in the tree } T_n\}.$$

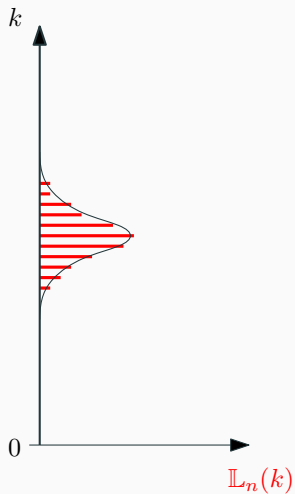
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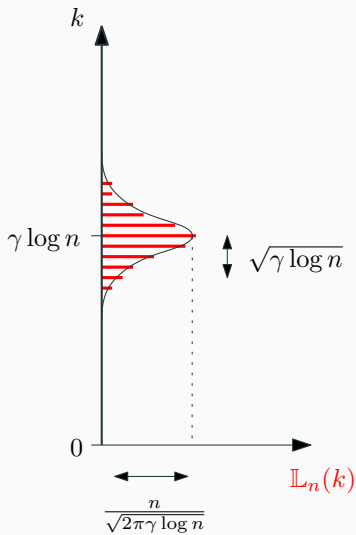
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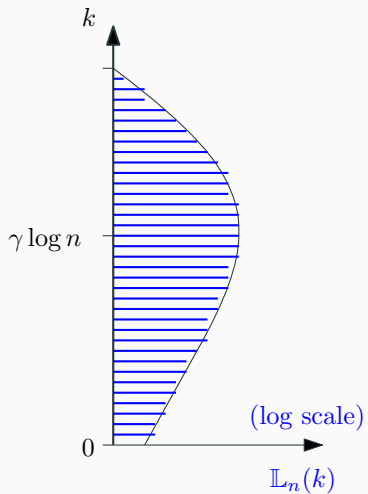
# THE PROFILE IS ALMOST SURELY ASYMPTOTICALLY GAUSSIAN



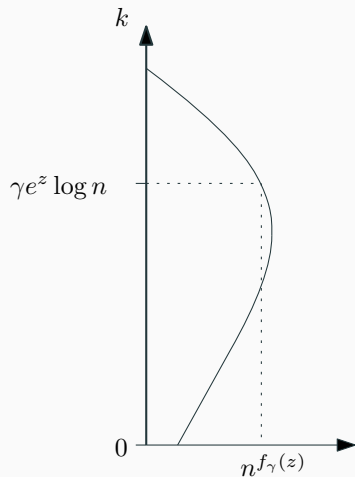
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# THE LOG OF THE PROFILE CONVERGES TO A FUNCTION

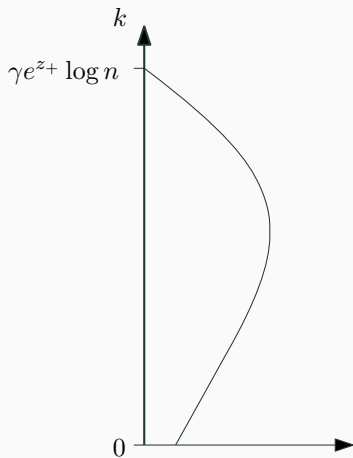


# THE LOG OF THE PROFILE CONVERGES TO A FUNCTION



$$\cdot f_\gamma(z) := 1 + \gamma(e^z - 1 - ze^z).$$

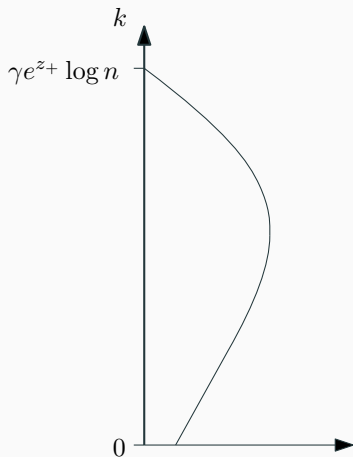
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- $z_+$  is the unique positive zero of  $f_\gamma$ .



## THE LOG OF THE PROFILE CONVERGES TO A FUNCTION



- $f_\gamma(z) := 1 + \gamma(e^z - 1 - ze^z)$ .
- $z_+$  is the unique positive zero of  $f_\gamma$ .
- Almost surely  
$$\text{ht}(T_n) \underset{n \rightarrow \infty}{\sim} \gamma \cdot e^{z_+} \log n.$$

## Theorem (S. 2019+)

Under the assumption  $W_n \simeq \text{cst} \cdot n^\gamma$  (+ additional condition) we almost surely have

- $\frac{\text{ht}(T_n)}{\log n} \xrightarrow{n \rightarrow \infty} \gamma \cdot e^{z_+}$ .
- $\mathbb{L}_n(k) \underset{n \rightarrow \infty}{=} \frac{n}{\sqrt{2\pi\gamma \log n}} \exp \left\{ -\frac{1}{2} \cdot \left( \frac{k - \gamma \log n}{\sqrt{\gamma \log n}} \right)^2 \right\} + O\left(\frac{n}{\log n}\right)$ ,
- for all  $z \in (z_-, z_+)$ ,  $\mathbb{L}_n(\lfloor \gamma e^z \log n \rfloor) = n^{f_\gamma(z) + o(1)}$ .