

A counterexample on the reconstruction of oligomorphic clones

Manuel Bodirsky, David Evans, **Michael Kompatscher**,
Michael Pinsker

michael@logic.at

Theory and Logic group
Technische Universität Wien

Homogeneous structures
BIRS - 13/11/2015

Reconstruction

Definition

We say a closed group $\Sigma \leq \text{Sym}(\omega)$ has **reconstruction** iff for every $\Sigma' \leq \text{Sym}(\omega)$:

$$\Sigma \cong \Sigma' \Rightarrow \Sigma \cong_{\mathcal{T}} \Sigma',$$

where $\cong_{\mathcal{T}}$ denotes topological isomorphism in the topology of pointwise convergence.

Reconstruction

Definition

We say a closed group $\Sigma \leq \text{Sym}(\omega)$ has **reconstruction** iff for every $\Sigma' \leq \text{Sym}(\omega)$:

$$\Sigma \cong \Sigma' \Rightarrow \Sigma \cong_T \Sigma',$$

where \cong_T denotes topological isomorphism in the topology of pointwise convergence.

We similarly define reconstruction for

- transformation monoids in ω^ω
- function clones in $\bigcup_{n \geq 1} \omega^{\omega^n}$.

Clones

$\mathcal{C} \subseteq \bigcup_{n \geq 1} A^{A^n}$ is called **function clone** if

- it contains all projections $\pi_i^n(x_1, \dots, x_n) = x_i$
- it is closed under composition
 $f(g_1(x_1, \dots, x_m), \dots, g_n(x_1, \dots, x_m))$

A map $\xi : \mathcal{C} \rightarrow \mathcal{D}$ is a **clone homomorphism** if

- ξ preserves arities
- $\xi(\pi_i^n) = \pi_i^n$
- it commutes with composition:
 $\xi(f(g_1, \dots, g_n)) = \xi(f)(\xi(g_1), \dots, \xi(g_n))$

$\bigcup_{n \geq 1} \omega^{\omega^n}$... topology of pointwise convergence.

ω -categorical structures

Ahlbrandt & Ziegler '86

Two countable ω -categorical structures \mathcal{A}, \mathcal{B} are **bi-interpretable** iff

$$\text{Aut}(\mathcal{A}) \cong_T \text{Aut}(\mathcal{B}).$$

- What information is contained in $\text{Aut}(\mathcal{A})$ as abstract group?
- Does every oligomorphic group have reconstruction?

ω -categorical structures

Ahlbrandt & Ziegler '86

Two countable ω -categorical structures \mathcal{A}, \mathcal{B} are **bi-interpretable** iff

$$\text{Aut}(\mathcal{A}) \cong_T \text{Aut}(\mathcal{B}).$$

- What information is contained in $\text{Aut}(\mathcal{A})$ as abstract group?
- Does every oligomorphic group have reconstruction?

Bodirsky & Pinsker '12

Two countable ω -categorical structures \mathcal{A}, \mathcal{B} are **primitive positive bi-interpretable** iff

$$\text{Pol}(\mathcal{A}) \cong_T \text{Pol}(\mathcal{B}).$$

- Does every oligomorphic clone have reconstruction?

Results

Evans & Hewitt '90

There is an ω -categorical structure \mathcal{A} , such that $\text{Aut}(\mathcal{A})$ has no reconstruction.

Results

Evans & Hewitt '90

There is an ω -categorical structure \mathcal{A} , such that $\text{Aut}(\mathcal{A})$ has no reconstruction.

Based on the group result:

Bodirsky, Evans, Pinsker & MK '15

There is an ω -categorical structure \mathcal{B} in *finite signature*, such that $\text{Aut}(\mathcal{B})$, $\text{End}(\mathcal{B})$ and $\text{Pol}(\mathcal{B})$ have no reconstruction.

Results

Evans & Hewitt '90

There is an ω -categorical structure \mathcal{A} , such that $\text{Aut}(\mathcal{A})$ has no reconstruction.

Based on the group result:

Bodirsky, Evans, Pinsker & MK '15

There is an ω -categorical structure \mathcal{B} in *finite signature*, such that $\text{Aut}(\mathcal{B})$, $\text{End}(\mathcal{B})$ and $\text{Pol}(\mathcal{B})$ have no reconstruction.

These are results in **ZFC**.

ZF(+DC) is consistent with: “Every isomorphism between closed subgroups of $\text{Sym}(\omega)$ is a homeomorphism.”

Proof outline

- 1 There are profinite $G_1, G_2 \leq \text{Sym}(\omega)$ such that $G_1 \cong G_2$ but $G_1 \not\cong_T G_2$.
- 2 Lift to oligomorphic $\Sigma_1, \Sigma_2 \leq \text{Sym}(\omega)$ such that $\Sigma_1 \cong \Sigma_2$ but $\Sigma_1 \not\cong_T \Sigma_2$.
- 3 The topological closures $\overline{\Sigma_1}$ and $\overline{\Sigma_2}$ in ω^ω are isomorphic, but not topologically isomorphic.
- 4 The clone closures $\text{Clo}(\overline{\Sigma_1}), \text{Clo}(\overline{\Sigma_2})$ are isomorphic, but not topologically isomorphic
- 5 $\text{Clo}(\overline{\Sigma_1})$ is the polymorphism clone of a structure in finite language.

The profinite group G

Fact (Witt '54)

There exists a separable profinite group G such that

- There is a finite central subgroup F such that $G = F \times E$ for an $E \leq G$.
- Every complement of a central subgroup is dense in G .

The profinite group G

Fact (Witt '54)

There exists a separable profinite group G such that

- There is a finite central subgroup F such that $G = F \times E$ for an $E \leq G$.
- Every complement of a central subgroup is dense in G .

Fix F and E . The profinite group G/F is abstractly isomorphic to E , so $G \cong F \times G/F$.

The profinite group G

Fact (Witt '54)

There exists a separable profinite group G such that

- There is a finite central subgroup F such that $G = F \times E$ for an $E \leq G$.
- Every complement of a central subgroup is dense in G .

Fix F and E . The profinite group G/F is abstractly isomorphic to E , so $G \cong F \times G/F$.

Every isomorphism $F \times G/F \rightarrow G$ maps the image of the compact group $1 \times G/F$ to a dense subgroup of G . Therefore $G \not\cong_T F \times G/F$.

Proof outline

- ✓ There are profinite $G_1, G_2 \leq \text{Sym}(\omega)$ such that $G_1 \cong G_2$ but $G_1 \not\cong_T G_2$.
- ② Lift to oligomorphic $\Sigma_1, \Sigma_2 \leq \text{Sym}(\omega)$ such that $\Sigma_1 \cong \Sigma_2$ but $\Sigma_1 \not\cong_T \Sigma_2$.
- ③ The topological closures $\overline{\Sigma_1}$ and $\overline{\Sigma_2}$ in ω^ω are isomorphic, but not topologically isomorphic.
- ④ The clone closures $\text{Clo}(\overline{\Sigma_1}), \text{Clo}(\overline{\Sigma_2})$ are isomorphic, but not topologically isomorphic
- ⑤ $\text{Clo}(\overline{\Sigma_1})$ is the polymorphism clone of a structure in finite language.

Encoding profinite groups with oligomorphic groups

Proposition (Cherlin, Hrushovski)

For every separable profinite group R there is an oligomorphic Σ_R :

- There is a continuous homomorphism $\nu : \Sigma_R \rightarrow R$
- The kernel Φ is the intersection of open subgroups of finite index in Σ_R

Encoding profinite groups with oligomorphic groups

Proposition (Cherlin, Hrushovski)

For every separable profinite group R there is an oligomorphic Σ_R :

- There is a continuous homomorphism $\nu : \Sigma_R \rightarrow R$
- The kernel Φ is the intersection of open subgroups of finite index in Σ_R

Wishful thinking

Use the profinite counterexample:

$$G \not\cong_T F \times G/F \Rightarrow \Sigma_G \not\cong_T \Sigma_{F \times G/F}$$

$$G \cong F \times G/F \Rightarrow \Sigma_G \cong \Sigma_{F \times G/F}$$

Encoding profinite groups with oligomorphic groups

Proposition (Cherlin, Hrushovski)

For every separable profinite group R there is an oligomorphic Σ_R :

- There is a continuous homomorphism $\nu : \Sigma_R \rightarrow R$
- The kernel Φ is the intersection of open subgroups of finite index in Σ_R

Wishful thinking

Use the profinite counterexample:

$$G \not\cong_T F \times G/F \Rightarrow \Sigma_G \not\cong_T \Sigma_{F \times G/F}$$

$$G \cong F \times G/F \Rightarrow \Sigma_G \cong \Sigma_{F \times G/F}$$

Problem: We do not know if $\Sigma_G \cong \Sigma_{F \times G/F}$ holds.

G/F and G as permutation groups

We know that $G \cong F \times G/F$, but $G \not\cong_T F \times G/F$.
How does this translate to permutation groups?

G/F and G as permutation groups

We know that $G \cong F \times G/F$, but $G \not\cong_T F \times G/F$.
How does this translate to permutation groups?

- G has base of clopen subgroups $(G_i)_{i \in \omega}$:
 $1 = \bigcap_{i \in \omega} G_i$ and $F = \bigcap_{i \geq 1} G_i$.

G/F and G as permutation groups

We know that $G \cong F \times G/F$, but $G \not\cong_T F \times G/F$.
How does this translate to permutation groups?

- G has base of clopen subgroups $(G_i)_{i \in \omega}$:
 $1 = \bigcap_{i \in \omega} G_i$ and $F = \bigcap_{i \geq 1} G_i$.
- $G \rightarrow \text{Sym}(\bigcup_{i \in \omega} G/G_i)$ is faithful and continuous.

G/F and G as permutation groups

We know that $G \cong F \times G/F$, but $G \not\cong_T F \times G/F$.
How does this translate to permutation groups?

- G has base of clopen subgroups $(G_i)_{i \in \omega}$:
 $1 = \bigcap_{i \in \omega} G_i$ and $F = \bigcap_{i \geq 1} G_i$.
- $G \rightarrow \text{Sym}(\bigcup_{i \in \omega} G/G_i)$ is faithful and continuous.
- $G/F \rightarrow \text{Sym}(\bigcup_{i \geq 1} G/G_i)$ is faithful and continuous.

G/F and G as permutation groups

We know that $G \cong F \times G/F$, but $G \not\cong_T F \times G/F$.
How does this translate to permutation groups?

- G has base of clopen subgroups $(G_i)_{i \in \omega}$:
 $1 = \bigcap_{i \in \omega} G_i$ and $F = \bigcap_{i \geq 1} G_i$.
- $G \rightarrow \text{Sym}(\bigcup_{i \in \omega} G/G_i)$ is faithful and continuous.
- $G/F \rightarrow \text{Sym}(\bigcup_{i \geq 1} G/G_i)$ is faithful and continuous.
- $F \times G/F \cong G$ gives us an action of G/F on G/G_0 .

G/F and G as permutation groups

We know that $G \cong F \times G/F$, but $G \not\cong_T F \times G/F$.

How does this translate to permutation groups?

- G has base of clopen subgroups $(G_i)_{i \in \omega}$:
 $1 = \bigcap_{i \in \omega} G_i$ and $F = \bigcap_{i \geq 1} G_i$.
- $G \rightarrow \text{Sym}(\bigcup_{i \in \omega} G/G_i)$ is faithful and continuous.
- $G/F \rightarrow \text{Sym}(\bigcup_{i \geq 1} G/G_i)$ is faithful and continuous.
- $F \times G/F \cong G$ gives us an action of G/F on G/G_0 .
- $G/F \rightarrow \text{Sym}(G/G_0 \cup \bigcup_{i \geq 1} G/G_i)$ is not continuous.
The closure of the image of G/F is isomorphic to G .

From $\Sigma_{G/F}$ to Γ

Similarly start with $\Sigma_{G/F}$, oligomorphic on A .

- Let $\Sigma_{G/F}$ act on G/G_0 via $\nu : \Sigma_{G/F} \rightarrow G/F$.
- The action $i : \Sigma_{G/F} \rightarrow \text{Sym}(A \cup G/G_0)$ is not continuous.
- Let Γ be the closure of $i(\Sigma_{G/F})$ in $\text{Sym}(A \cup G/G_0)$.

From $\Sigma_{G/F}$ to Γ

Similarly start with $\Sigma_{G/F}$, oligomorphic on A .

- Let $\Sigma_{G/F}$ act on G/G_0 via $\nu : \Sigma_{G/F} \rightarrow G/F$.
- The action $i : \Sigma_{G/F} \rightarrow \text{Sym}(A \cup G/G_0)$ is not continuous.
- Let Γ be the closure of $i(\Sigma_{G/F})$ in $\text{Sym}(A \cup G/G_0)$.

Result for groups

- Γ is oligomorphic on $A \cup G/G_0$,
- $\Gamma \cong i(\Sigma_{G/F}) \times F$,
- G is the quotient of Γ and the intersection of open subgroups with finite index.

Conclusion: $\Gamma \cong \Sigma_{G/F} \times F$ but $\Gamma \not\cong_T \Sigma_{G/F} \times F$

Proof outline

- ✓ There are profinite $G_1, G_2 \leq \text{Sym}(\omega)$ such that $G_1 \cong G_2$ but $G_1 \not\cong_T G_2$.
- ✓ Lift to oligomorphic $\Sigma_1, \Sigma_2 \leq \text{Sym}(\omega)$ such that $\Sigma_1 \cong \Sigma_2$ but $\Sigma_1 \not\cong_T \Sigma_2$.
- ③ The topological closures $\overline{\Sigma_1}$ and $\overline{\Sigma_2}$ in ω^ω are isomorphic, but not topologically isomorphic.
- ④ The clone closures $\text{Clo}(\overline{\Sigma_1}), \text{Clo}(\overline{\Sigma_2})$ are isomorphic, but not topologically isomorphic
- ⑤ $\text{Clo}(\overline{\Sigma_1})$ is the polymorphism clone of a structure in finite language.

Lifting to the monoid closure

We know that $\bar{\Gamma} \not\cong_T F \times \overline{\Sigma_{G/F}}$.

Lifting to the monoid closure

We know that $\bar{\Gamma} \not\cong_T F \times \overline{\Sigma_{G/F}}$.

To show: $\bar{\Gamma}$ and $F \times \overline{\Sigma_{G/F}}$ are isomorphic as abstract monoids.

Lifting to the monoid closure

We know that $\bar{\Gamma} \not\cong_T F \times \overline{\Sigma_{G/F}}$.

To show: $\bar{\Gamma}$ and $F \times \overline{\Sigma_{G/F}}$ are isomorphic as abstract monoids.

Problem: How to lift a *non-continuous* isomorphism $\Gamma \cong \Sigma_{G/F} \times F$ to the topological closure?!

Lifting to the monoid closure

We know that $\bar{\Gamma} \not\cong_T F \times \overline{\Sigma_{G/F}}$.

To show: $\bar{\Gamma}$ and $F \times \overline{\Sigma_{G/F}}$ are isomorphic as abstract monoids.

Problem: How to lift a *non-continuous* isomorphism $\Gamma \cong \Sigma_{G/F} \times F$ to the topological closure?!

Lemma

The encoding homomorphism $\nu : \Sigma_{G/F} \rightarrow G/F$ extends to a continuous monoid homomorphism

$$\nu : \overline{\Sigma_{G/F}} \rightarrow G/F \text{ with kernel } \bar{\Phi}.$$

Lifting to the monoid closure

Define via $\nu : \overline{\Sigma_{G/F}} \rightarrow G/F$ an action of $\overline{\Sigma_{G/F}}$ on G/G_0 .
The combined action

$$i : \overline{\Sigma_{G/F}} \rightarrow (A \cup G/G_0)^{(A \cup G/G_0)}$$

is not continuous. The topological closure of $i(\overline{\Sigma_{G/F}})$ is $\overline{\Gamma}$.

Lifting to the monoid closure

Define via $\nu : \overline{\Sigma_{G/F}} \rightarrow G/F$ an action of $\overline{\Sigma_{G/F}}$ on G/G_0 .
The combined action

$$i : \overline{\Sigma_{G/F}} \rightarrow (A \cup G/G_0)^{(A \cup G/G_0)}$$

is not continuous. The topological closure of $i(\overline{\Sigma_{G/F}})$ is $\overline{\Gamma}$.

Lifting to the monoid closure

Define via $\nu : \overline{\Sigma_{G/F}} \rightarrow G/F$ an action of $\overline{\Sigma_{G/F}}$ on G/G_0 .
The combined action

$$i : \overline{\Sigma_{G/F}} \rightarrow (A \cup G/G_0)^{(A \cup G/G_0)}$$

is not continuous. The topological closure of $i(\overline{\Sigma_{G/F}})$ is $\overline{\Gamma}$.

Similar to the group case:

- $\overline{\Gamma}$ is an oligomorphic monoid,
- $\overline{\Gamma} \cong i(\overline{\Sigma_{G/F}}) \times F$,

Lifting to the monoid closure

Define via $\nu : \overline{\Sigma_{G/F}} \rightarrow G/F$ an action of $\overline{\Sigma_{G/F}}$ on G/G_0 .
The combined action

$$i : \overline{\Sigma_{G/F}} \rightarrow (A \cup G/G_0)^{(A \cup G/G_0)}$$

is not continuous. The topological closure of $i(\overline{\Sigma_{G/F}})$ is $\overline{\Gamma}$.

Similar to the group case:

- $\overline{\Gamma}$ is an oligomorphic monoid,
- $\overline{\Gamma} \cong i(\overline{\Sigma_{G/F}}) \times F$,

Result for monoids

$\overline{\Gamma}$ and $\overline{\Sigma_{G/F}} \times F$ are isomorphic, but not topologically isomorphic.

Proof outline

- ✓ There are profinite $G_1, G_2 \leq \text{Sym}(\omega)$ such that $G_1 \cong G_2$ but $G_1 \not\cong_T G_2$.
- ✓ Lift to oligomorphic $\Sigma_1, \Sigma_2 \leq \text{Sym}(\omega)$ such that $\Sigma_1 \cong \Sigma_2$ but $\Sigma_1 \not\cong_T \Sigma_2$.
- ✓ The topological closures $\overline{\Sigma_1}$ and $\overline{\Sigma_2}$ in ω^ω are isomorphic, but not topologically isomorphic.
- ④ The clone closures $\text{Clo}(\overline{\Sigma_1}), \text{Clo}(\overline{\Sigma_2})$ are isomorphic, but not topologically isomorphic
- ⑤ $\text{Clo}(\overline{\Sigma_1})$ is the polymorphism clone of a structure in finite language.

The clone closure

Let $\text{Clo}(\Lambda)$ denote the smallest clone containing the transformation monoid Λ . So

$$f(x_1, \dots, x_n) \in \text{Clo}(\Lambda) \Leftrightarrow f(x_1, \dots, x_n) = g(x_i)$$

for $g \in \Lambda$, $1 \leq i \leq n$.

The clone closure

Let $\text{Clo}(\Lambda)$ denote the smallest clone containing the transformation monoid Λ . So

$$f(x_1, \dots, x_n) \in \text{Clo}(\Lambda) \Leftrightarrow f(x_1, \dots, x_n) = g(x_i)$$

for $g \in \Lambda$, $1 \leq i \leq n$.

Observation

Let $I : \Lambda \rightarrow \Delta$ be a monoid isomorphism. If I and I^{-1} send constants to constants, it has a natural extension to a clone isomorphism $\text{Clo}(\Lambda) \rightarrow \text{Clo}(\Delta)$.

The clone closure

Let $\text{Clo}(\Lambda)$ denote the smallest clone containing the transformation monoid Λ . So

$$f(x_1, \dots, x_n) \in \text{Clo}(\Lambda) \Leftrightarrow f(x_1, \dots, x_n) = g(x_i)$$

for $g \in \Lambda$, $1 \leq i \leq n$.

Observation

Let $I : \Lambda \rightarrow \Delta$ be a monoid isomorphism. If I and I^{-1} send constants to constants, it has a natural extension to a clone isomorphism $\text{Clo}(\Lambda) \rightarrow \text{Clo}(\Delta)$.

Result for clones

The clones $\text{Clo}(\bar{\Gamma})$ and $\text{Clo}(\overline{\Sigma_{G/F}} \times F)$ are isomorphic but not topologically isomorphic.

Proof outline

- ✓ There are profinite $G_1, G_2 \leq \text{Sym}(\omega)$ such that $G_1 \cong G_2$ but $G_1 \not\cong_T G_2$.
- ✓ Lift to oligomorphic $\Sigma_1, \Sigma_2 \leq \text{Sym}(\omega)$ such that $\Sigma_1 \cong \Sigma_2$ but $\Sigma_1 \not\cong_T \Sigma_2$.
- ✓ The topological closures $\overline{\Sigma_1}$ and $\overline{\Sigma_2}$ in ω^ω are isomorphic, but not topologically isomorphic.
- ✓ The clone closures $\text{Clo}(\overline{\Sigma_1}), \text{Clo}(\overline{\Sigma_2})$ are isomorphic, but not topologically isomorphic
- ⑤ $\text{Clo}(\overline{\Sigma_1})$ is the polymorphism clone of a structure with finite signature.

Finite language

Proposition (Hrushovski)

Let \mathcal{A} be ω -categorical. Then there is an ω -categorical \mathcal{B} in a finite language including a predicate P , such that:

- the domain of \mathcal{A} is equal to $P^{\mathcal{B}}$
- the definable relations in \mathcal{A} are exactly the definable relations of \mathcal{B} on P

Finite language

Proposition (Hrushovski)

Let \mathcal{A} be ω -categorical. Then there is an ω -categorical \mathcal{B} in a finite language including a predicate P , such that:

- the domain of \mathcal{A} is equal to $P^{\mathcal{B}}$
- the definable relations in \mathcal{A} are exactly the definable relations of \mathcal{B} on P

Idea

Let \mathcal{A} be such that $\Sigma_{G/F} = \text{Aut}(\mathcal{A})$. Take \mathcal{B} given by Hrushovski's construction.

Make sure that $\text{Aut}(\mathcal{B})$ still “encodes” G/F and $\text{End}(\mathcal{B}) = \overline{\text{Aut}(\mathcal{B})}$. Then $F \times \text{Aut}(\mathcal{B})$ and $F \times \text{End}(\mathcal{B})$ have no reconstruction.

Thank you!