

Hypercube percolation

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Hypercube percolation

Study random subgraph of hypercube $\{0, 1\}^n$, with bonds

$$\{\{x, y\} : \exists! i \text{ with } x_i \neq y_i\}.$$

Make bonds $\{x, y\}$ independently

occupied with probability p ,

vacant with probability $1 - p$,

where $p \in [0, 1]$ is percolation parameter.

Erdős and Spencer (1979): for $p = 1/2 + \varepsilon/(2n)$, random graph connected with probability $e^{-e^{-\varepsilon}(1+o(1))}$.

Our goal: Study percolation phase transition as $n \rightarrow \infty$.

Inspiration: Erdős-Rényi random graph

Erdős-Rényi random graph is random subgraph of complete graph on n vertices where each of $\binom{n}{2}$ edges is occupied with probab. p .

Phase transition: (Erdős and Rényi (60))

For $p = (1 + \varepsilon)/n$, largest connected component $|\mathcal{C}_{\max}|$ has size

(a) $\Theta_{\mathbb{P}}(\log n)$ for $\varepsilon < 0$;

(b) $\Theta_{\mathbb{P}}(n)$ for $\varepsilon > 0$;

Scaling window: (Bollobás (84) and Łuczak (90))

For $p = (1 + \varepsilon)/n$, largest component size is $\Theta_{\mathbb{P}}(n^{2/3})$ with expected cluster size $\Theta(n^{1/3})$ whenever $|\varepsilon| \leq \Lambda n^{-1/3}$.

Barely sub- and supercritical regimes:

(a) $|\mathcal{C}_{\max}| = 2\varepsilon n(1 + o_{\mathbb{P}}(1))$ when $\varepsilon \gg n^{-1/3}$, $\varepsilon = o(1)$;

(b) $|\mathcal{C}_{\max}| = 2\varepsilon^{-2} \log(\varepsilon^3 n)(1 + o_{\mathbb{P}}(1))$ when $\varepsilon \ll -n^{-1/3}$, $\varepsilon = o(1)$.

BCHSS approach to criticality

Recall $\mathbb{E}_p |\mathcal{C}(0)| = \Theta(n^{1/3})$ for Erdős-Rényi random graph and $p = 1/n$, where n is volume graph.

BCHSS05a,05b,06: Define critical threshold $p_c(\{0, 1\}^n)$ as solution

$$\chi(p) = 2^{n/3},$$

where $\chi(p) = \mathbb{E}_p |\mathcal{C}(0)|$ is expected cluster size on hypercube.

Challenge: Prove that $p_c(\{0, 1\}^n)$ really is critical value, by investigating connected component sizes above and below $p_c(\{0, 1\}^n)$.

(Sub-)Critical results

Theorem 1 (Subcritical clusters) (BCHSS (05a), (05b)).

For $p = p_c(1 + \varepsilon)$, and uniformly in $\varepsilon \leq 0$, as $n \rightarrow \infty$,

$$\chi(p) = \frac{O(1)}{|\varepsilon| + 2^{-n/3}},$$

$$\mathbb{P}_p\left(\chi^2(p) \leq |\mathcal{C}_{\max}| \leq 2\chi^2(p) \log(2^n/\chi^3(p))\right) \geq 1 - \log(2^n/\chi^3(p))^{-3/2}.$$

Theorem 2 (Scaling window) (BCHSS (05a), (05b)).

For $p = p_c(1 + \varepsilon)$, with $|\varepsilon| \leq \Lambda 2^{-n/3}$, there exists $b_1 = b_1(\Lambda) > 0$ s.t.

$$\mathbb{P}_p\left(\omega^{-1}2^{2n/3} \leq |\mathcal{C}_{\max}| \leq \omega 2^{2n/3}\right) \geq 1 - \frac{b_1}{\omega}.$$

vdHHe (11): $|\mathcal{C}_{\max}|2^{-2n/3}$ is tight and non-degenerate:

hallmark of critical behavior.

Percolation on high-dimensional tori

Results work on general high-dimensional tori $\mathbb{T}_{r,d}$, subject to a finite graph triangle condition.

Define $p_c(\mathbb{T}_{r,d}; \lambda)$ such that $\mathbb{E}_{p_c} |\mathcal{C}(0)| = \lambda r^{d/3}$, and let

$$\tau_p(x) = \mathbb{P}_p(0 \longleftrightarrow x).$$

Then, finite-graph triangle condition states that

$$\nabla(p_c(\mathbb{T}_{r,d}; \lambda)) = \sum_{x,y} \tau_{p_c}(x) \tau_{p_c}(y-x) \tau_{p_c}(y) = 1 + o(1) + O(\lambda^3).$$

Known for

- ▷ Percolation on nearest-neighbor torus $\mathbb{T}_{r,d}$ for d large;
- ▷ Percolation on spread-out torus and $d > 6$;
- ▷ Hypercube percolation.

[Beware, not obvious that $p_c(\mathbb{T}_{r,d}; \lambda)$ close to $p_c(\mathbb{Z}^d)$ for high-d tori...]

Supercritical results

Theorem 3 (Supercritical clusters) (vdH+Nachmias (12)).

For $p = p_c(1 + \varepsilon)$, with $\varepsilon = o(1), \varepsilon \gg 2^{-n/3}$,

$$\frac{|\mathcal{C}_{\max}|}{2\varepsilon 2^n} \xrightarrow{\mathbb{P}} 1,$$

while, with $\mathcal{C}_{(2)}$ denoting **second largest cluster**,

$$\frac{|\mathcal{C}_{(2)}|}{\varepsilon 2^n} \xrightarrow{\mathbb{P}} 0,$$

$$\mathbb{E}_p |\mathcal{C}(0)| = 4\varepsilon^2 2^n (1 + o(1)).$$



Percolation phase transitions hypercube and complete graph alike:
 2ε is asymptotic survival probability of branching process with
Poisson($1 + \varepsilon$)-offspring distribution.

Proof (sub)critical regime BCHSS

(a) Differential inequalities and **finite-graph triangle condition** (as in Aizenman+Newman 84, Barsky+Aizenman 91);

(b) **Lace expansion for two-point function** (as in Hara+Slade 90) to relate finite-graph triangle condition to **random walk condition**;

(c) **Analysis (a)-(b)** proving bounds on expected and largest cluster.

▷ Part (b) is only one that uses **lace expansion**.

(Later explain how **lace expansion** can be avoided for hypercube.)

General supercritical setting

Main result follows from **general result** applying to many **high-dimensional transitive graphs** of large degree m .

Applies under **three conditions** on transitive base graph on which percolation is performed:

- (1) **degree m tends to infinity** with volume graph;
- (2) $[(m - 1)p_c]^{T_{\text{mix}}} = 1 + o(1)$, where T_{mix} is **mixing time non-backtracing random walk**;
- (3) **finite graph NBW triangle condition**.

Proof 1: supercritical cluster tails

Theorem 4 (Supercritical cluster tails). Let $p = p_c(1 + \varepsilon)$ with $\varepsilon \gg V^{-1/3}$. Then, for $k_0 = \varepsilon^{-2}(\varepsilon^3 V)^\alpha$ for any $\alpha \in (0, 1/3)$,

$$\mathbb{P}_p(|\mathcal{C}(0)| \geq k_0) = 2\varepsilon(1 + o(1)).$$

Further, with

$$Z_{\geq k} = |\{v : |\mathcal{C}(v)| \geq k\}|,$$

denoting number of vertices in large clusters,

$$\frac{Z_{\geq k_0}}{2\varepsilon 2^n} \xrightarrow{\mathbb{P}} 1.$$

Proof: Differential inequalities (as in BA91, BCHSS05a), with careful analysis of constants involved.

Proof 2: non-backtracking random walk

Let $\mathbb{P}_p(0 \overset{[a,b]}{\longleftrightarrow} x)$ denote probability that

0 is connected to x with **shortest path of length** $\in [a, b]$.

We bound $\mathbb{P}_p(0 \overset{[r,2r]}{\longleftrightarrow} x)$ by a constant in x for $r \gg 1/\varepsilon$ by comparing percolation paths to **non-backtracking random walk** (i.e., random walk conditioned not to reverse immediately).

In particular, with $V = 2^n$,

$$\mathbb{P}_p(0 \overset{[r,2r]}{\longleftrightarrow} x) \leq \frac{1 + o(1)}{V} \sum_x \mathbb{P}_p(0 \overset{[r,2r]}{\longleftrightarrow} x) \leq (1 + o(1)) \mathbb{E}_p |B_0(2r)| / V,$$

where $B_0(r)$ is ball radius r in graph metric on G_p .

This makes performing **complicated sums** relatively easy.

Proof 3: large clusters share boundary

For $x, y \in \{0, 1\}^n$ and ℓ , let

$$S_\ell(x, y) = |\{(u, u') \in E : x \overset{\ell}{\leftrightarrow} u, y \overset{\ell}{\leftrightarrow} u'\}|.$$

Pair x, y is **good** when $|\mathcal{C}(x)|, |\mathcal{C}(y)| \geq K\varepsilon^{-2}$, and

$$S_{2r+r_0}(x, y) \geq \alpha_n(nV) \left(\mathbb{E}|B_0(r_0)| / (2\varepsilon) \right)^2 V^{-2},$$

where $\alpha_n \rightarrow 0$ very slowly. Write $P_{r,r_0,K}$ for number of good pairs.

Theorem 5 (Most large clusters share many boundary edges).

Assume $\varepsilon^3 2^n \rightarrow \infty, \varepsilon \leq n^{-2}$. Take $r = M/\varepsilon, r_0 = \frac{1}{2}\varepsilon^{-1} \log(\varepsilon^3 2^n)[1 - o(1)]$.

Then, there exist $K, M \rightarrow \infty$ s.t.

$$\frac{P_{r,r_0,K}}{(2\varepsilon V)^2} \xrightarrow{\mathbb{P}} 1.$$

Proof 4: improved sprinkling

Take $p_2 = \delta\varepsilon/n$ for our sprinkling probability, where $\delta > 0$ is small.
Let p_1 satisfy

$$p = p_1 + (1 - p_1)p_2,$$

Given G_{p_1} , construct **auxiliary simple graph** $H = (\mathcal{V}, \mathcal{E})$ with

$$\mathcal{V} = \{x \in G_{p_1} : |\mathcal{C}(x)| \geq K\varepsilon^{-2}\}, \quad \mathcal{E} = \{(x, y) \in \mathcal{V}^2 \text{ good pair}\},$$

$$\mathbb{P}_{p_1}(P_{r, r_0, K} \geq (1 - \alpha)4\varepsilon^2 2^{2n}, |\mathcal{V}| \in [(2 - \theta)\varepsilon, (2 + \theta)\varepsilon]) \geq 1 - \delta.$$

Claim: Probability there exists partition $\mathcal{V} = M_1 \uplus M_2$ with $|M_1| \geq \theta|\mathcal{V}|$ and $|M_2| \geq \theta|\mathcal{V}|$ s.t. **sprinkled edges** do not connect M_1 to M_2 is **small**.

Step 1: Number of partitions is at most $2^{3K^{-1}\varepsilon^3V}$.

Step 2: Given such partition, number of edges (u, u') between $\mathcal{C}(M_1)$ and $\mathcal{C}(M_2)$ is at least $\alpha_n^3 \varepsilon^2 n 2^n$ for some large K . [Cheat]

Proof 4: improved sprinkling (Cont.)

- ▷ Recall that $p_2 = \delta\varepsilon/n$ is sprinkling probability, with $\delta > 0$ small.
- ▷ Recall that number of edges (u, u') between $\mathcal{C}(M_1)$ and $\mathcal{C}(M_2)$ is at least $\alpha_n^3 \varepsilon^2 n 2^n$ for some large K .

Conclusion: Probability sprinkling fails is at most

$$2^{3K-1} \varepsilon^{3V} \left(1 - \delta\varepsilon/n\right)^{\alpha_n K^{-2} \varepsilon^2 n 2^n} \leq 2^{3K-1} \varepsilon^{3V} e^{-\alpha_n^3 \delta \varepsilon^3 2^n},$$

which converges to zero rapidly when $\varepsilon^3 2^n \gg 1$ and $K \alpha_n^3 \gg 1$.

Can achieve this by letting $\alpha_n \rightarrow 0$ and $K \rightarrow \infty$ sufficiently slowly. For example, $K = (\varepsilon^3 2^n)^{1/4}$, $\alpha_n = (\varepsilon^3 2^n)^{-1/13}$ will do.

Proof 5: Unlacing hypercube percolation

For hypercube, we aim to show that

$$[(n-1)p_c]^{T_{\text{mix}}} = 1 + o(1),$$

where T_{mix} is mixing time non-backtracing random walk, which equals $T_{\text{mix}} = (1 + o(1))n \log(n)$ (vdH-Fitzner 14).

This amounts to showing $\mathbb{E}_p |\mathcal{C}(0)| \gg 2^{n/3}$ for $p = 1/(n-1) + K/n^3$ with $K > 0$ large.

Allows us to identify phase transition hypercube without relying on lace expansion.

Lemma. For $p = (1 + 5/(2n^2) + K/n^3)/(n-1)$ with K sufficiently large, as long as $\mathbb{E}_p |B(k)| \leq 2^{n/2}$, there exists $c > 0$ s.t.

$$\mathbb{E}_p |\partial B(k)| \geq [1 + c/n^3] \mathbb{E}_p |\partial B(k-1)|.$$

High-dimensional tori

Theorem 5. (Heydenreich-vdH (07-11)) For nearest-neighbour percolation in sufficiently high dimensions, or for sufficiently spread-out percolation in $d > 6$,

$$p_c(\mathbb{Z}^d) = p_c(\mathbb{T}_{r,d}; \lambda) + O(r^{-d/3}).$$

As a result, all results from BCHSS05a apply.

Proof uses subtle coupling of clusters on torus and \mathbb{Z}^d , similar in spirit to that by Benjamini-Schramm (96), that shows that

$$|\mathcal{C}_{\mathbb{T}}(0)| \preceq |\mathcal{C}_{\mathbb{Z}}(0)|,$$

together with an accompanying lower bound.



Progress nearest-neighbor model

Theorem 6. (Fitzner-vdH (14?)) Nearest-neighbour percolation displays mean-field critical exponents for $d \geq 12$.



Proof by Hara-Slade was never published, so project improves transparency in high-dimensional percolation.

Mathematica code can be downloaded from website Robert Fitzner, for everyone to check and play with.

Proof uses

- (a) Non-backtracking lace expansion (NoBLE) taking more of interaction explicitly into account;
- (b) sharp bounds on NoBLE coefficients;
- (c) analysis of NoBLE assuming explicit numerical conditions;
- (c) a computer-assisted proof that verifies these conditions.

Open problems and extension

- (a) Study limit in probability of **largest cluster** in subcritical phase on hypercube.
- (b) Study limit in probability of **second largest component** in supercritical phase on hypercube, and prove **discrete duality principle**.
- (c) Prove CLT for **giant component**.
- (d) Show that **scaling limit** of largest clusters is same (modulo trivial scalings) as that for ERRG in Aldous (97).
- (e) Extension **percolation on torus to nearest-neighbor percolation** with $d \geq 12$.

Literature hypercube percolation

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