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New inequalities for the  
von Neumann entropy (?)

Andreas Winter (U. Bristol & NUS)

with Noah Linden, CMP 259, 129-138 (2005)

Ben Ibrison, PhD Thesis, Bristol (2007)  
[unpublished]

# Outline

I. Entropy + Inequalities

II. "Laws of Info. Th." for 2 & 3 parties

III. 4 parties: constrained inequality

IV. Quest for an unconstr. inequality...

V. Open problems

# I. Entropy + Inequalities

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$S(\rho) = -\text{Tr} \rho \log \rho$  von Neumann entropy  
for state  $\rho \geq 0$ ,  $\text{Tr} \rho = 1$

$$= -\sum_j \lambda_j \log \lambda_j \quad \text{for } \rho = \sum_j \lambda_j |e_j\rangle\langle e_j|$$

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Universal relations between entropies of  
multipartite systems?

I.e., for density operator  $\rho_{ABC\dots Z}$  on Hilbert space  $A \otimes B \otimes C \dots \otimes Z$ :

Consider marginals  $\rho_A = \text{Tr}_{B\dots Z} \rho$   
 $\rho_B = \text{Tr}_{AC\dots Z} \rho$   
 $\rho_{AB} = \text{Tr}_{C\dots Z} \rho$   
etc.

Entropies  $S(A) = S(\rho_A)$   
 $S(B) = S(\rho_B)$   
 $S(AB) = S(\rho_{AB})$   
etc.

For  $n$ -party state:  $2^n$  values, one for each subset  $\mathcal{X}$

Relations:

$$S(\mathcal{X}) \geq 0 \text{ and } S(\emptyset) = 0$$

(0)

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$$I(A:B) := S(A) + S(B) - S(AB) \geq 0$$

(1)

$$I(A:C|B) := S(AB) + S(BC) - S(ABC) - S(B) \geq 0$$

(2)

( "strong subadditivity" — Lieb/Ruskai JRP '73 )

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[but not in general  $S(\mathcal{X}) \geq S(\mathcal{Y})$   
for  $\mathcal{X} \supset \mathcal{Y}$  !]

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[ $\mathcal{X} \cap \mathcal{Y} = \emptyset \Rightarrow S(\mathcal{X}) + S(\mathcal{Y}) - S(\mathcal{X} \cup \mathcal{Y}) \geq 0$ ]

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[ $S(\mathcal{X}) + S(\mathcal{Y}) - S(\mathcal{X} \cup \mathcal{Y}) - S(\mathcal{X} \cap \mathcal{Y}) \geq 0$ ]

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...but there is more — purely quantum mechanically:

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- (i) If  $\rho = |\psi\rangle\langle\psi|$  is pure (i.e.  $S(\rho) = 0$ ), then  
 $S(\mathcal{X}) = S(\overline{\mathcal{X}})$  [since  $\rho_{\mathcal{X}}$  and  $\rho_{\overline{\mathcal{X}}}$  have the same spectrum]
- (ii) Any ( $n$ -party) state  $\rho$  is the marginal of an ( $n+1$ -party) pure state  $|\psi\rangle\langle\psi|$ .

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Hence:

$$S(AB) + S(A) - S(B) \geq 0 \tag{3}$$

$$[\mathcal{X} \cap \mathcal{Y} = \emptyset \Rightarrow S(\mathcal{X} \cup \mathcal{Y}) + S(\mathcal{X}) - S(\mathcal{Y}) \geq 0]$$

$$S(AB) + S(BC) - S(A) - S(C) \geq 0 \tag{4}$$

$$[S(\mathcal{X}) + S(\mathcal{Y}) - S(\mathcal{X} \setminus \mathcal{Y}) - S(\mathcal{Y} \setminus \mathcal{X}) \geq 0]$$

“weak monotonicity”

I.e., all entropy vectors  $[S(\mathcal{X}) : \emptyset \neq \mathcal{X} \subset \{1, \dots, n\}]$  <sup>16</sup>

$\uparrow$   
 $\mathbb{R}^{2^n - 1}$

are contained in the convex  
cone  $\Sigma_n$ , defined by

set of  $\mathcal{X}$   
all:  $\Sigma_n^*$

(Positivity)

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Pippenger IEEE-IT '03:  $\overline{\Sigma_n^*}$  also a convex cone, contained in  $\Sigma_n$ : i.e. with  $\vec{v}, \vec{w} \in \overline{\Sigma_n^*} \forall \lambda, \mu \geq 0$ , contains  $\lambda \vec{v} + \mu \vec{w}$ .

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# Questions:

• Is  $\overline{\Sigma}_n^* = \Sigma_n$ ? or  $\neq$ ?

• How to describe  $\overline{\Sigma}_n^*$ ?

→ Try to discover all inequalities!

[cf. Pippenger, "What are the laws of Information Theory?", 1986]

Classical case:  $\overline{\Delta}_n^* \subset \Delta_n$  ← Shannon cone  
 R.W. Yeung; Z. Zhang; R. Dougherty et al.; ...

•  $\Delta_n^*$  not closed  
 •  $\overline{\Delta}_n^* \neq \Delta_n$ :  $\exists$  non-Shannon-type ineq's!

## II. $n=1, 2$ and $3$ parties

$n=1$   $\Sigma_1^* = \Sigma_1$  : all non-negative values allowed.



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$n=1$

$$\Sigma_1^* = \Sigma_1$$

: all non-negative values allowed.

$n=2$

$$\Sigma_2^* = \Sigma_2$$

: all non-negative values of  $S(A), S(B), S(AB)$  that obey

$$S(AB) \leq S(A) + S(B)$$
$$S(A) \leq S(AB) + S(B)$$
$$S(B) \leq S(AB) + S(A)$$

are allowed.

$n=3$

Pippenger IEEE-IT '03 computed the  
 external rays of  $\Sigma_3$  : spanned by

A	B	C	AB	AC	BC	ABC
1	1	0	0	1	1	0
1	0	0	1	1	0	1
1	1	1	1	1	1	1
1	1	1	2	2	2	1

(up to permutations)

$n=3$

Pippenger IEEE-IT '03 computed the  
extremal rays of  $\Sigma_3$ : spanned by

A	B	C	AB	AC	BC	ABC
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1	0	0	1	1	0	1
1	1	1	1	1	1	1
1	1	1	2	2	2	1

(up to permutations)

... sufficient to find states realizing them!

$n=3$   
cont'd

Useful to introduce purifying system  
D, e.t.  $S(D) = S(ABD)$ ,  $S(AD) = S(BC)$ , etc.

A	B	C	D	AB	AC	BC
1	1	0	0	0	1	1
1	1	1	1	1	1	1
1	1	1	1	2	2	2

$n=3$   
cont'd

Useful to introduce purifying system  
 $D$ , s.t.  $S(D) = S(AB), S(AD) = S(BC)$ , etc.

A	B	C	D	AB	AC	BC	
1	1	0	0	0	1	1	(i)
1	1	1	1	1	1	1	(ii)
1	1	1	1	2	2	2	(iii)

States:

- (i)  $|EPR\rangle_{AB} |0\rangle_C |0\rangle_D = \frac{1}{\sqrt{2}} (|0000\rangle + |1100\rangle)$
- (ii)  $|GHZ\rangle_{ABCD} = \frac{1}{\sqrt{2}} (|0000\rangle + |1111\rangle)$
- (iii)  $\frac{1}{3} \sum_{a,b \in \mathbb{Z}_3} |a\rangle_A |b\rangle_B |a+b\rangle_C |a+2b\rangle_D =: |\Psi\rangle_{ABCD}$   
"max. entangled state"

### III. 4 parties...

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...are similarly attacked via extremal rays of  $\Sigma_4$

[ Linden / Manera / Masser / Popescu /  
Roberts / Schumacher / Smolin /  
Thapliyal, unpublished '03 ]

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With purifying system  $E$ , there are 76 rays.

Up to permutation symmetry, only 8 classes,  
as follows...

	A	B	C	D	E	AB	AC	AD	AE	BC	BD	BE	CD	CE	DE
					ABCD				BCD			ACD		ABD	ABC
(i)	1	1	0	0	0	0	1	1	1	1	1	1	0	0	0
(ii)	1	1	1	1	0	1	1	1	1	1	1	1	1	1	1
(iii)	1	1	1	1	0	2	2	2	1	2	2	1	2	1	1
(iv)	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
(v)	2	1	1	1	1	3	3	3	3	2	2	2	2	2	2
(vi)	1	1	2	2	2	2	3	3	3	3	3	3	2	2	2
(vii)	3	3	2	2	2	4	3	3	3	3	3	3	3	4	4
(viii)	3	3	3	3	2	4	4	4	5	4	4	5	6	5	5



(i), (ii), (iii) we know already : realized by

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$|EPR\rangle_{AB} |000\rangle_{CDE}$

$|GHZ\rangle_{ABCDE} |0\rangle_E$

$|M\rangle_{ABCD} |0\rangle_E$



(i), (ii), (iii) we know already : realized by

$$\begin{aligned}
 &|EPR\rangle_{AB} |000\rangle_{CDE} \\
 &|\Psi\rangle_{ABCD} |0\rangle_E \\
 &|\Psi\rangle_{ABCD} |0\rangle_E
 \end{aligned}$$

(iv) : realized by  $|\Psi\rangle_{ABCDE}$

(v), (vi) : realized by code or secret sharing states

$$\text{e.g. } |\Psi_5\rangle = \frac{1}{\sqrt{2}} [ |0\rangle_A |0\rangle_{L} + |1\rangle_A |1\rangle_{L} ]_{ABCDE}$$

Logical 0, 1 in 5-qubit code

$$|\Psi_6\rangle = \frac{1}{\sqrt{27}} \sum_{a,b,c \in \mathbb{Z}_3} |a\rangle_A |b\rangle_B |a+b\rangle_C |a+c\rangle_C |a+b\rangle_D |a+c\rangle_D |a+b\rangle_E |a+c\rangle_E$$

(vii), (viii) : ???

Thm. (Linden/AW '05) If  $(ABC D)$  satisfies

$$I(A:B|C) = 0$$

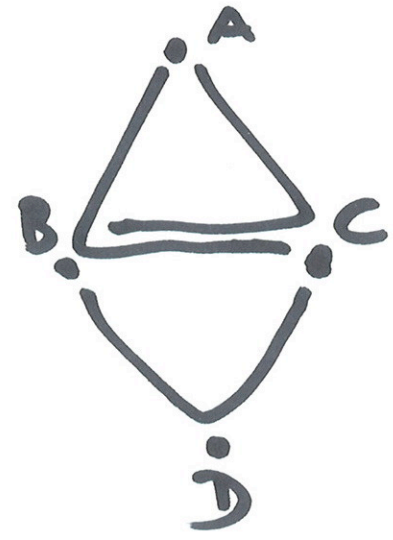
$$I(A:C|B) = 0$$

$$I(B:C|D) = 0$$

(i.e. SSA is saturated),

then

$$I(A:D) \geq I(A:BC).$$



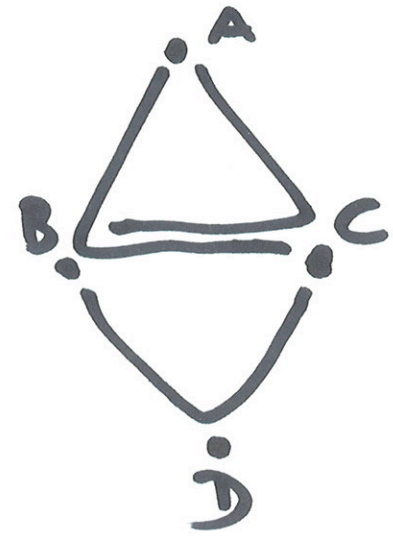
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then  $I(A:D) \geq I(A:BC)$ .

Proof uses characterization of "E" case of SSA [Hayden/Massa/Pete/AW, CMP '03] as "quantum Markov chains" & exploiting the ensuing classical information ...

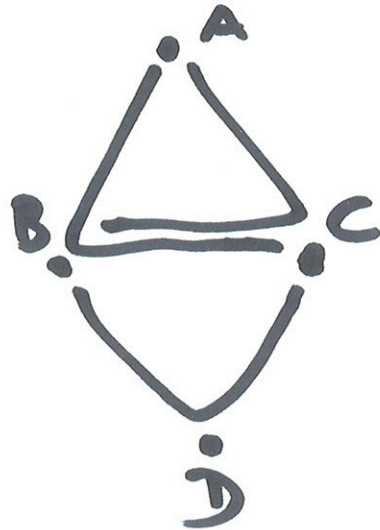
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Both rays (vii) & (viii) contradict the Thm.

$\Rightarrow \Sigma_q^* \neq \Sigma_q$  (but leaves open  $\overline{\Sigma_q^*} \stackrel{?}{=} \Sigma_q$ )

## IV. An unconstrained inequality?

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New (constrained) inequality cannot be derived from "standard" ones: it excludes part of a face of  $\Sigma_4$ .

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Here some attempts to find one ( $\rightarrow$  Ben Ibinson, PhD thesis, Bristol 2007)



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Hence, there are more extremal rays to  $\overline{\Sigma_4^*}$ .

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So, perhaps there exist  $\kappa_i \geq 0$  s.t. for all states,

$$\kappa_1 I(A:C|B) + \kappa_2 I(B:C|A) + \kappa_3 I(A:B|D)$$

$$+ [I(C:D) - I(C:AB)] \geq 0$$

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$$+ [I(C:D) - I(C:AB)] \geq 0$$

... However, difficult to test since  $k_i$  unknown

[see also Ibinson/Linden/Alw, CMP '08]

(6) Young - Zhang inequality

Young / Zhang, IEEE-IT '98 showed that the following is true for classical densities:

$$2I(A:B) + 3I(A:B|C) + I(A:B|D) \leq I(C:D) + I(C:AB)$$

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For quantum states:

- violated by rays (vii) & (viii)

- no counterexamples known (neither in special states nor in numerical search)

- YZ proof (via custom-made conditional indep. structure) doesn't work, however...



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- \* ... but as of now have only weak indications what they might be
- \* Strongest candidate: Yang-Zhang inequality
- \* Interesting to note: All known parts of the cones  $\overline{\Sigma_n^*}$  spanned by stabiliser states.