

An introduction to matroid theory

- The definition and some basic examples
- Elementary operations on matroids
- Some important theorems

<http://www.math.LSU.edu/~oxley/survey4.pdf>

A fundamental example

$$A = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 & 5 & 6 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \end{matrix} & \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix} \end{matrix}$$

$$E = \{1, 2, 3, 4, 5, 6\} \quad \begin{matrix} \text{column} \\ \text{labels} \end{matrix}$$

\mathcal{I} : lin. independent sets of columns

A_3 : view A over $GF(3)$

A_2 : view A over $GF(2)$

$$\mathcal{I}_3 = \{X \subseteq E: |X| \leq 3 \text{ and} \\ X \neq \{1, 2, 4\}, \{1, 3, 5\}, \{2, 3, 6\}\}$$

\mathcal{I}_2 : same except $X \neq \{4, 5, 6\}$

$$A = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 & 5 & 6 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \end{matrix} & \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix} \end{matrix}$$

$$\mathcal{I}_3 = \{X \subseteq E : |X| \leq 3 \text{ and } X \neq \{1, 2, 4\}, \\ \{1, 3, 5\}, \{2, 3, 6\}\}$$

\mathcal{I}_2 : same except $X \neq \{4, 5, 6\}$

(E, \mathcal{I}_3) and (E, \mathcal{I}_2) are matroids.

ground set independent sets

(I1) \mathcal{I} is non-empty.

(I2) Every subset of an indep't set is indep't.

(I3) If $X \subseteq E$, all max'l indep't subsets of X have the same size.

(I1) \mathcal{Y} is non-empty.

(I2) Every subset of an indep't set is indep't.

(I3) If $X \subseteq E$, all max'l indep't subsets of X have the same size.

A matroid $M = (E, \mathcal{Y})$ is a finite set E and a collection \mathcal{Y} of subsets of E satisfying

(I1) - (I3).

Examples Vector matroids of matrices

$M[A_3]$

ternary $(GF(3))$

$M[A_2]$

binary $(GF(2))$

(I1) \mathcal{I} is non-empty.

(I2) Every subset of an indept set is indept.

(I3) If $X \in E$, all max'l indept subsets of X have the same size.

maximal independent sets : bases

dependent set : not indept

minimal dependent sets : circuits

rank $r(X)$ of X : size of a max'l indept subset of X .

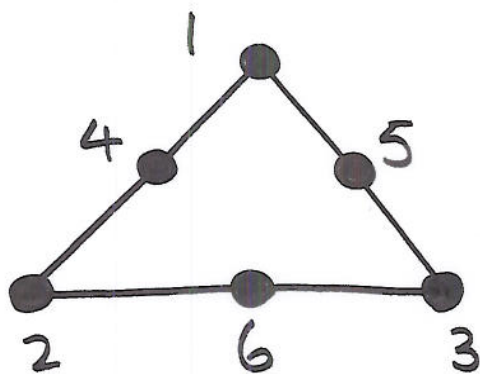
Example

$$A = \begin{bmatrix} 1 & 2 & 3 & 4 & 5 & 6 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}$$

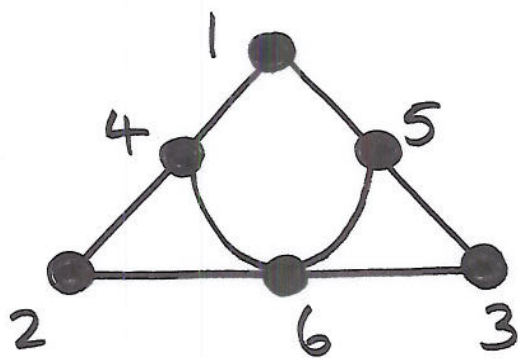
$$r(\{4,5,6\})$$

$$= \begin{cases} 3 & \text{in } M[A_3] \\ 2 & \text{in } M[A_2] \end{cases}$$

A geometric example



- two distinct lines have at most one common point.



This gives a matroid.

$$\mathcal{I}_{\text{top}} = \{X \subseteq \{1, 2, \dots, 6\} : |X| \leq 3 \text{ and } X \neq \{1, 2, 4\}, \{1, 3, 5\}, \{2, 3, 6\}\}$$

$$\mathcal{I}_{\text{bottom}} = \text{same except } X \neq \{4, 5, 6\}$$

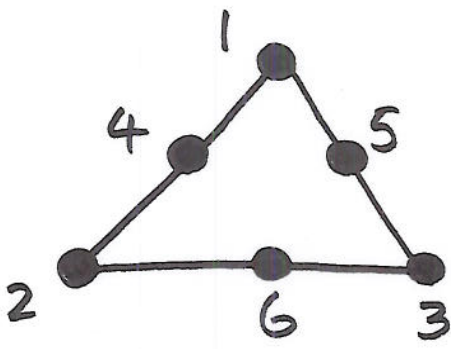
E : a set of points in the plane

Set of lines : sets of ≥ 2 points

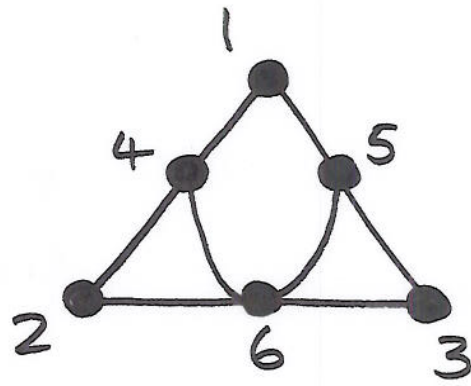
lines have at most one common point.

Dependent sets:

- 3 collinear points
- 4 coplanar points



M_t



M_b

For a field F , a matroid M is F -representable if $M \cong M[A]$ for some matrix A over F .

$GF(2)$ -rep.

binary

$GF(3)$ -rep.

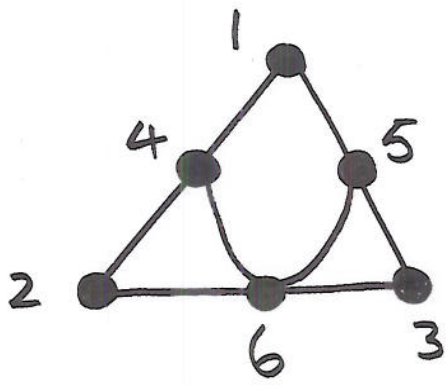
ternary

M_t is ternary

$$\begin{bmatrix} 1 & 2 & 3 & 4 & 5 & 6 \\ 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}_{GF(3)}$$

M_b is binary

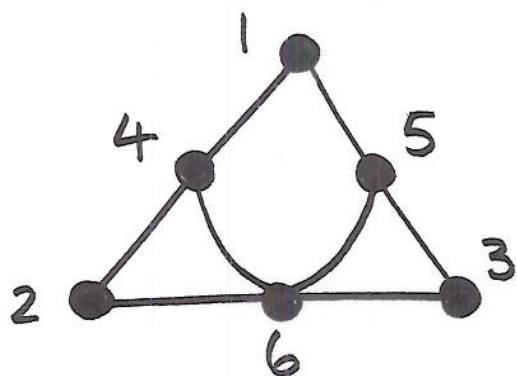
$$\begin{bmatrix} 1 & 2 & 3 & 4 & 5 & 6 \\ 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}_{GF(2)}$$



M_b

M_b is binary.

Is M_b ternary?



M_b is binary.

Is M_b ternary?

$$A_2 = \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}$$

binary rep'n

$$D_3 = \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & -1 \end{bmatrix}$$

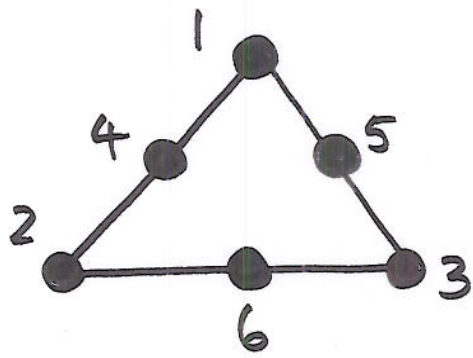
ternary rep'n.

M_b is regular (representable over every field).

$M_b = M[D_3]$ where D_3 is interpreted over the field given.

D_3 is a totally unimodular matrix

- all its subdeterminants calculated over \mathbb{R} are in $\{0, 1, -1\}$.



M_t

M_t is ternary.

Is M_t binary?

Try to build a binary rep'n.

$\{1, 2, 3\}$ is
a basis :

$$\begin{bmatrix} 1 & 2 & 3 & 4 & 5 & 6 \\ 1 & 0 & 0 & & & \\ 0 & 1 & 0 & & & \\ 0 & 0 & 1 & & & \end{bmatrix}$$

$\{1, 2, 4\}$ is
a circuit :

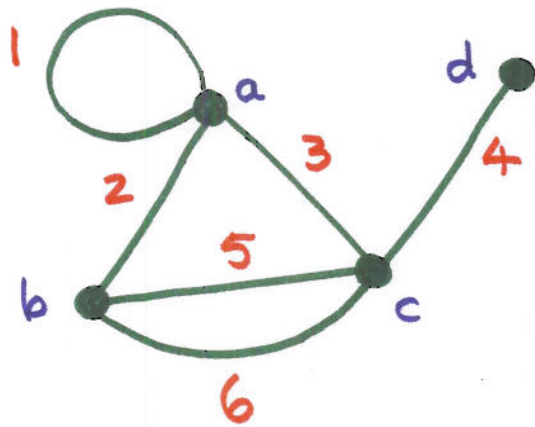
$$\begin{bmatrix} 1 & 2 & 3 & 4 & 5 & 6 \\ 1 & 0 & 0 & 1 & & \\ 0 & 1 & 0 & 1 & & \\ 0 & 0 & 1 & 0 & & \end{bmatrix}$$

$\{1, 3, 5\}$ and
 $\{2, 3, 6\}$ are
circuits :

$$\begin{bmatrix} 1 & 2 & 3 & 4 & 5 & 6 \\ 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}_{GF(2)}$$

But $M_t \neq M[A_2]$, so M_t is
non-binary.

Matroids from graphs



G

	1	2	3	4	5	6
a	0	1	1	0	0	0
b	0	1	0	0	1	1
c	0	0	1	1	1	1
d	0	0	0	1	0	0

Vertex-edge incidence matrix over $GF(2)$

A_G

$M(G)$, cycle matroid of G , is $M[A_G]$.

circuits of $M(G)$: cycles of G

Some questions

Q1. Is every matroid F -rep. for some field F ?

Q2. If not, what can be said about $\frac{\# \text{ } F\text{-rep. } n\text{-element matroids}}{\# \text{ } n\text{-element matroids}}$?

Q3. For a specific field $GF(q)$, which matroids are $GF(q)$ -representable?

Which graphs are planar?

Minors of matroids

$$A_3 = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 & 5 & 6 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \end{matrix} & \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix} \end{matrix}$$

$M[A_3] \setminus 1$ (delete 1)

is represented by $\begin{matrix} & \begin{matrix} 2 & 3 & 4 & 5 & 6 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \end{matrix} & \begin{bmatrix} 0 & 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 1 \end{bmatrix} \end{matrix}$

For an element e of a matroid $M = (E, \mathcal{I})$,

the deletion $M \setminus e$ is the matroid $(E - \{e\}, \mathcal{I}')$ where

$$\mathcal{I}' = \{X \subseteq E - \{e\} : X \in \mathcal{I}\}.$$

- Easy to check this is a matroid.

$$A_3 = \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}$$

What is $M[A_3] / 1$ (contract 1)?

$$A_3 = \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}$$

$M[A_3] / 1$
is rep. by

$$\begin{bmatrix} 2 & 3 & 4 & 5 & 6 \\ 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 1 \end{bmatrix}$$

What about contracting 4 instead?

Pivot on
the circled
entry.

$$\begin{bmatrix} 1 & 2 & 3 & 4 & 5 & 6 \\ 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}$$

Row 2 - Row 1

$$\begin{bmatrix} 1 & 2 & 3 & 4 & 5 & 6 \\ -1 & 0 & 0 & 1 & 1 & 0 \\ -1 & 1 & 0 & 0 & -1 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}$$

$M[A_3] / 4$
is rep. by

$$\begin{bmatrix} 1 & 2 & 3 & 5 & 6 \\ -1 & 1 & 0 & -1 & 1 \\ 0 & 0 & 1 & 1 & 1 \end{bmatrix}$$

Contraction

If $\{e\}$ is a loop (one-element circuit) of a matroid M , then

$$M/e = M \setminus e.$$

If $\{e\}$ is independent and $M = (E, \mathcal{I})$,

then $M/e = (E - \{e\}, \mathcal{I}'')$ (M contract e)

where $\mathcal{I}'' = \{X \subseteq E - \{e\} : X \cup \{e\} \in \mathcal{I}\}$.

- M/e is a matroid.

- $M \setminus e \setminus f = M \setminus f \setminus e$

$$M/e/f = M/f/e$$

$$M \setminus e/f = M/f \setminus e$$

- $M \setminus X / Y$ is a minor of M ,

for disjoint subsets X and Y of E .

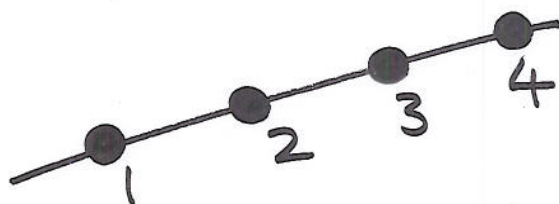
Lemma Every minor of an F -rep. matroid is F -rep.

excluded minors for F -rep.: matroids that are not F -rep. but every deletion and every contraction is F -rep.

Example

$$E = \{1, 2, 3, 4\}$$

$$\mathcal{Y} = \{X \subseteq E : |X| \leq 2\}$$



$$(E, \mathcal{Y}) = U_{2,4}$$

$U_{2,4}$ is not binary:

a 2-dimensional vector space over

$GF(2)$ has $2^2 - 1$ non-zero vectors.

Theorem (Tutte 1958) A matroid

is binary iff it has no $U_{2,4}$ -minor.

Binary matroids have many other characterizations.

$$X \Delta Y = (X - Y) \cup (Y - X)$$

Theorem The following are equivalent for a matroid M :

- (i) M is binary.
- (ii) For all distinct circuits C_1 and C_2 , the set $C_1 \Delta C_2$ contains a circuit.
- (iii) For all distinct circuits C_1 and C_2 , $C_1 \Delta C_2$ is a disjoint union of circuits.

Duality

$$A_3 = \begin{bmatrix} 1 & 2 & 3 & 4 & 5 & 6 \\ 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}$$

- a standard rep. matrix for $M[A_3]$
- a standard-form generator matrix for the $[6,3]$ -code over $GF(3)$, the code C is the vector space generated by the rows of A_3
- $A_3 = [I_3 | D]$

$[-D^T | I_{6-3}]$ is a parity-check matrix for C , a generator matrix for C^\perp , the dual code.

- The vector matroid of $[-D^T | I_{6-3}]$

i.e. of

$$\begin{bmatrix} 1 & 2 & 3 & 4 & 5 & 6 \\ -1 & -1 & 0 & 1 & 0 & 0 \\ -1 & 0 & -1 & 0 & 1 & 0 \\ 0 & -1 & -1 & 0 & 0 & 1 \end{bmatrix}$$

is the dual matroid of $M[A_3]$.

For a matroid M with ground set E and with \mathcal{B} as its set of bases, the dual matroid M^* has ground set E and its set of bases is $\{E-B: B \in \mathcal{B}\}$.

Lemma $[I_r | D]$ represents an n -element matroid M over a field F iff $[-D^T | I_{n-r}]$ represents M^* over F .

- $(M^*)^* = M$

- $M / T = (M^* \setminus T)^*$

- $M \setminus T = (M^* / T)^*$

- A matroid N is a minor of M iff N^* is a minor of M^* .

- M is an excluded minor for F -rep. $\Leftrightarrow M^*$ is an excluded minor for F -rep.

Terminology

M^*
circuit
basis

M
cocircuit
cobasis

A set is **spanning** if it contains a basis.

- (i) X is a basis of $M \Leftrightarrow E-X$ is a basis of M^* .
- (ii) X is spanning in $M \Leftrightarrow X$ contains a basis of M .
 $\Leftrightarrow E-X$ is contained in a basis of M^* .
 $\Leftrightarrow E-X$ is indept in M^* .
- (iii) X is non-spanning in $M \Leftrightarrow E-X$ is dept in M^* .

(iii) X is non-spanning in $M \iff E-X$ is dependent in M^* .

(iv) X is a max'l non-spanning set in $M \iff E-X$ is a minimal dep't set in M^* .

$\iff E-X$ is a circuit of M^*

$\iff E-X$ is a cocircuit of M .

maximal non-spanning set: hyperplane

Lemma The hyperplanes of M are the complements of its cocircuits.

Corollary A circuit and a cocircuit in a matroid cannot have exactly one common element.

Circuit - cocircuit intersections

Theorem The following are equivalent for a matroid M .

- (i) M is binary.
- (ii) $|C \cap C^*|$ is even for every circuit C and cocircuit C^* .
- (iii) $|C \cap C^*| \neq 3$ for every circuit C and cocircuit C^* .

Cocircuits and weights

Lemma Let A be a matrix over a field F , and let $M = M[A]$. The cocircuits of M coincide with the minimal non-empty supports of vectors from the row space of A .

- If A is a generator matrix of a code C , then the minimum weight of C is the minimum cocircuit size of $M[A]$.

Uniform matroids

Let $E = \{1, 2, \dots, n\}$ $0 \leq r \leq n$

$U_{r,n} = (E, \mathcal{I})$ where $\mathcal{I} = \{X \subseteq E : |X| \leq r\}$.

$$\underline{U_{r,n}^* = U_{n-r,n}}$$

- If $U_{r,n}$ is $GF(q)$ -rep., we have an $r \times n$ matrix A over $GF(q)$ such that every r columns are lin. indep't.

The $[n,r]$ -code over $GF(q)$ generated by A is a maximum-distance-separable (MDS) code.

Main conjecture for MDS codes

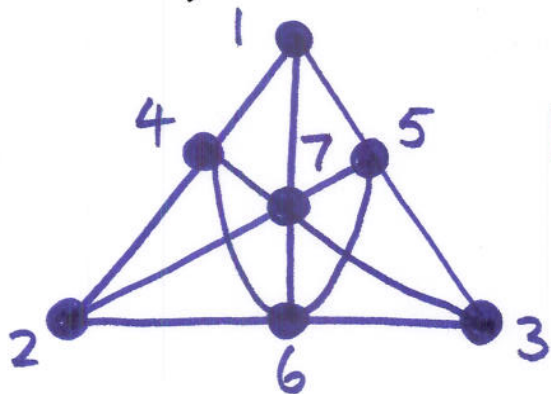
Conjecture For all prime powers q , $U_{r, q+2}$ is an excluded minor for the class of $GF(q)$ -rep. matroids

$\left\{ \begin{array}{l} \text{for all } r \text{ in } \{2, 3, \dots, q\} \text{ for } q \text{ odd;} \\ \text{for all } r \text{ in } \{2, 3, \dots, q\} - \{3, q-1\} \text{ for } q \text{ even.} \end{array} \right.$

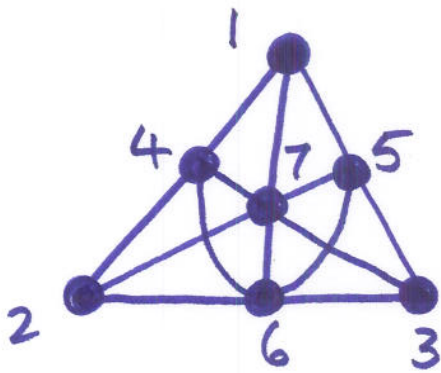
- Holds for $r \leq 5$.
- Holds for $q \leq 27$.

Excluded minors for ternary matroids

- $U_{2,5}$ and its dual, $U_{3,5}$



F_7 , the Fano matroid,
and its dual, F_7^*



F_7

$$\begin{bmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ 1 & 0 & 0 & 1 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 \end{bmatrix}_{GF(2)}$$

The dual code is the binary Hamming code $\text{Ham}(3, 2)$; a rep'n for F_7^* is a generator matrix for $\text{Ham}(3, 2)$.

Theorem (Reid; Bixby; Seymour)
1971 1979 1979

A matroid is ternary iff it has no minor isomorphic to any of

$U_{2,5}$, $U_{3,5}$, F_7 , or F_7^* .

Rota's Conjecture (1970)

For all finite fields $GF(q)$, there are only finitely many excluded minors for the class of $GF(q)$ -rep. matroids.

$$q = 2$$

$$U_{2,4}$$

$$q = 3$$

$$U_{2,5}, U_{3,5}, F_7, F_7^*$$

Theorem (Geelen, Gerards, Kapoor 2000)

There are exactly seven excluded minors for the class of $GF(4)$ -rep matroids.

$$q \geq 5$$

OPEN

- For all infinite fields F , there are infinitely many excluded minors for F -rep.

Regular matroids

(representable over every field)

Theorem (Tutte)

The following are equivalent for a matroid M .

- (i) M is regular.
- (ii) M is binary and ternary.
- (iii) M is binary and F -rep. for some field of characteristic not 2.

Class

binary

ternary

regular

{ = binary \wedge ternary }
Tutte (1958)

Excluded minors

$U_{2,4}$

$U_{2,5}, U_{3,5}, F_7, F_7^*$

minor-minimal members
of $\{U_{2,4}, U_{2,5}, U_{3,5}, F_7, F_7^*\}$
= $\{U_{2,4}, F_7, F_7^*\}$

Non-representable matroids

Proposition For a fixed field F ,

non-iso. F -rep. n -elt matroids $\rightarrow 0$

non-iso. n -element matroids

as $n \rightarrow \infty$.

Conjecture (Mayhew, Newman, Welsh, Whittle, 2009)

non-iso. rep. n -elt matroids $\rightarrow 0$

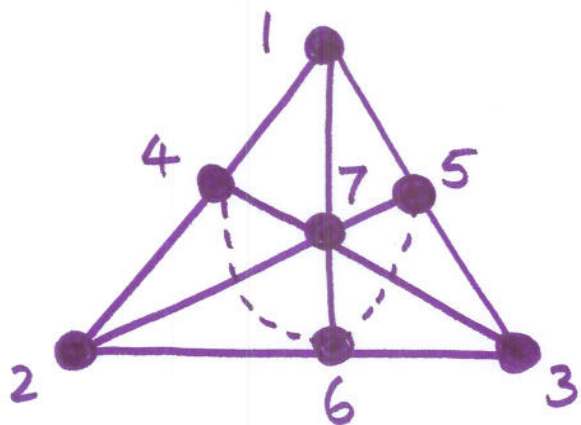
non-iso. n -element matroids

as $n \rightarrow \infty$

(rep: representable over some field)

Proposition Every matroid with at most 7 elements is representable; so is every matroid with 8 elements unless its rank is 4.

Example



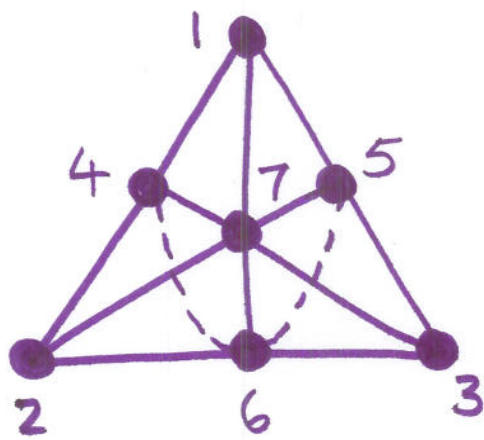
F_7 (Fano)

The dotted line encodes the equation

$$1+1 = 0$$

Proposition Every matroid with at most 7 elements is representable; so is every matroid with 8 elements unless its rank is 4.

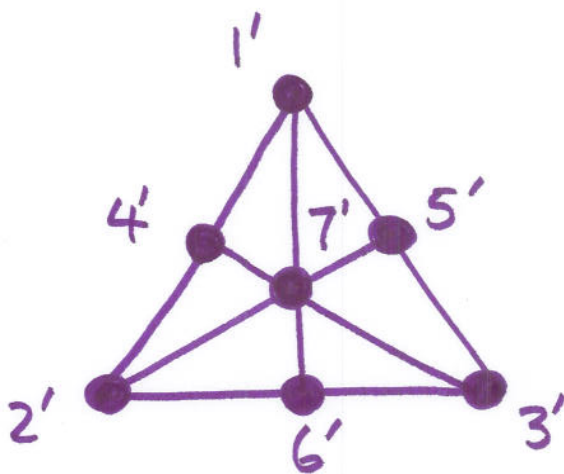
Example



F_7 (Fano)

$\{4, 5, 6\}$ is a circuit

$\Rightarrow 1+1 = 0$



F_7^- (non-Fano)

$\{4', 5', 6'\}$ is not a circuit

$\Rightarrow 1+1 \neq 0$

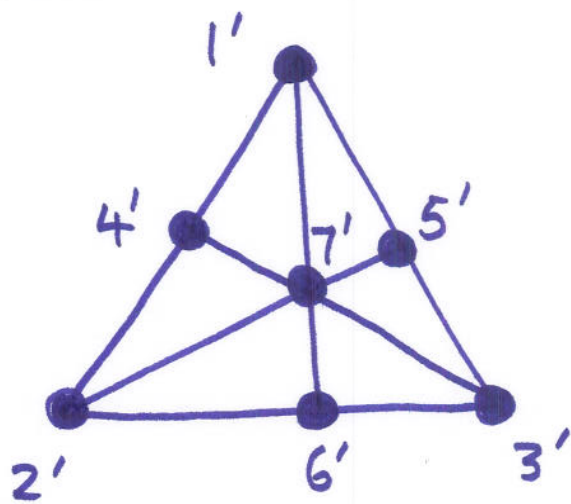
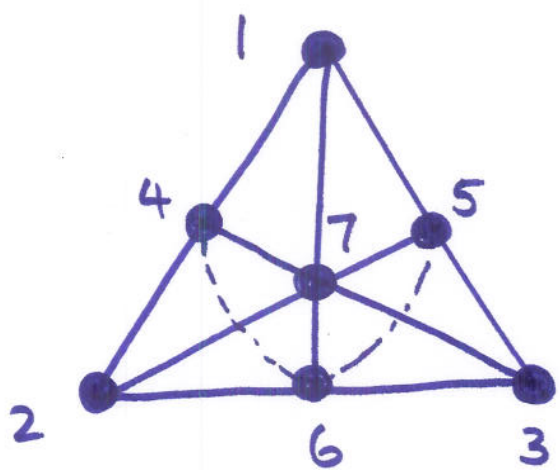
$F_7 \oplus F_7^-$ is not rep'ble.

This matroid has ground set $E(F_7) \cup E(F_7^-)$.

Its indept sets: $\{I_1 \cup I_2 : I_1 \in \mathcal{I}(F_7), I_2 \in \mathcal{I}(F_7^-)\}$

- $F_7 \oplus F_7^-$, the direct sum of F_7 and F_7^- , has 14 elements and $\text{rank} = 3 + 3 = 6$.

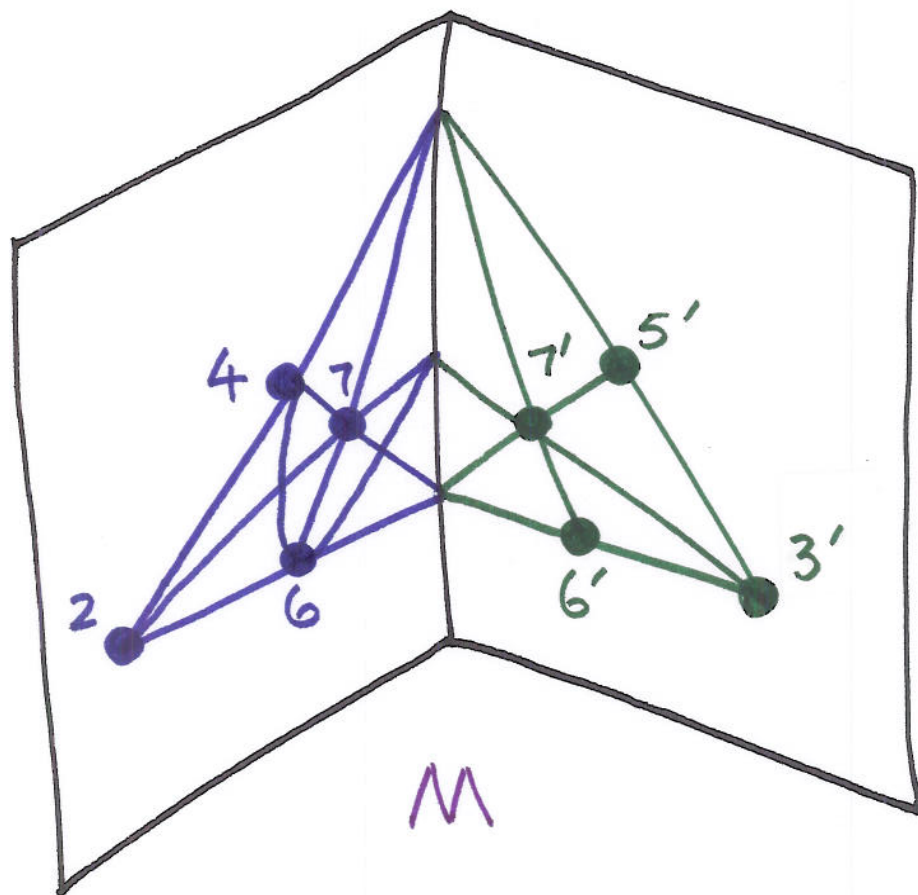
A different way to put F_7 and F_7^- together: in rank 3



$$E = \{1, 2, \dots, 7\} \cup \{1', 2', \dots, 7'\}$$

$$\mathcal{Y} = \{X \subseteq E : |X| \leq 3, \text{ elements of } X \text{ are not collinear}\}$$

Another way to join F_7 and F_7^-
in rank 4



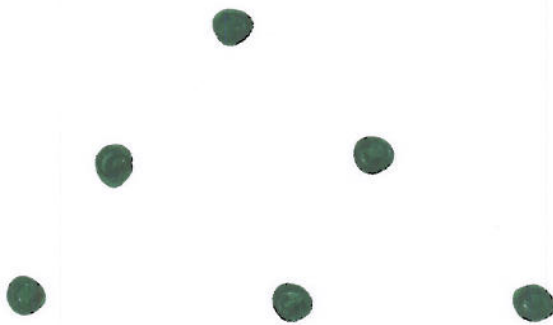
- 8 elements, rank 4

- $M / 2 \cong F_7^-$
 - $M / 3' \cong F_7$
- To contract 2, project from 2, so 4, 7, and 6 plug the holes in F_7^- .

$U_{r,n}$: n -element matroid

- every r -element set is a basis

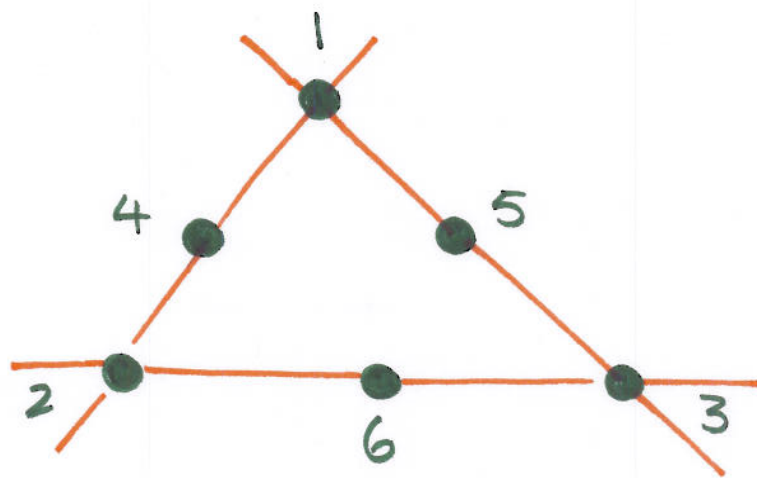
$U_{3,6}$



6 points
in the
plane, no
three
collinear.

$U_{r,n}$: n -element matroid

- every r -element set is a basis



Three orange lines correspond to circuits and hyperplanes: circuit-hyperplanes.

In general, take a subset \mathcal{X} of the bases of $U_{r,n}$ so that no two members of \mathcal{X} have more than $r-2$ common elements.

\mathcal{X} -sets are non-bases - they become circuit-hyperplanes.

This gives a sparse paving matroid

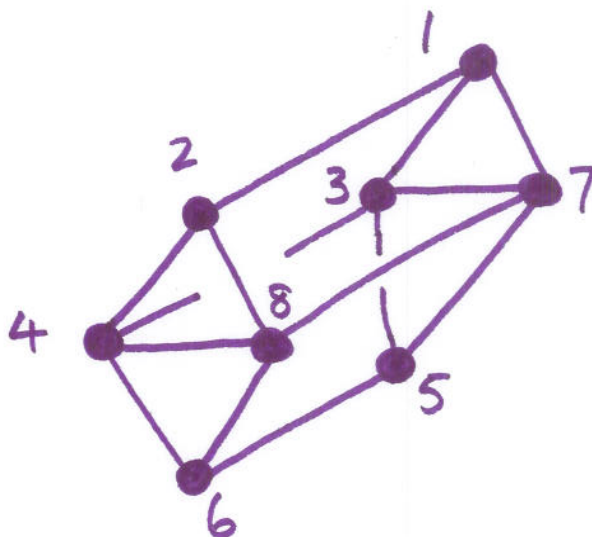
Example Start with $U_{4,8}$.

Let $\mathcal{X} = \{ \{1,2,3,4\}, \{3,4,5,6\}, \{5,6,7,8\}, \{1,2,7,8\}, \{3,4,7,8\} \}$

Matroid V_8 , the Vámos matroid

$E = \{1, 2, \dots, 8\}$

bases : all 4-element sets except those in \mathcal{X} .



This matroid is not representable.

Rank axioms

Lemma Let E be a finite set. A

function $r: 2^E \rightarrow \mathbb{Z}^+ \cup \{0\}$

is the rank function of a matroid on E

iff

(R1) If $X \subseteq E$, then $r(X) \leq |X|$.

(R2) If $X \subseteq Y \subseteq E$, then $r(X) \leq r(Y)$.

(increasing)

(R3) If $X, Y \subseteq E$, then

(submodular)

$$r(X) + r(Y) \geq r(X \cup Y) + r(X \cap Y)$$

Independent sets: $|X| = r(X)$

(R1) If $X \subseteq E$, then $r(X) \leq |X|$.

(R2) If $X \subseteq Y \subseteq E$, then $r(X) \leq r(Y)$.

(R3) If $X, Y \subseteq E$, then

$$r(X) + r(Y) \geq r(X \cup Y) + r(X \cap Y).$$

Weaken (R1) to

$$(R1)' \quad r(\emptyset) = 0.$$

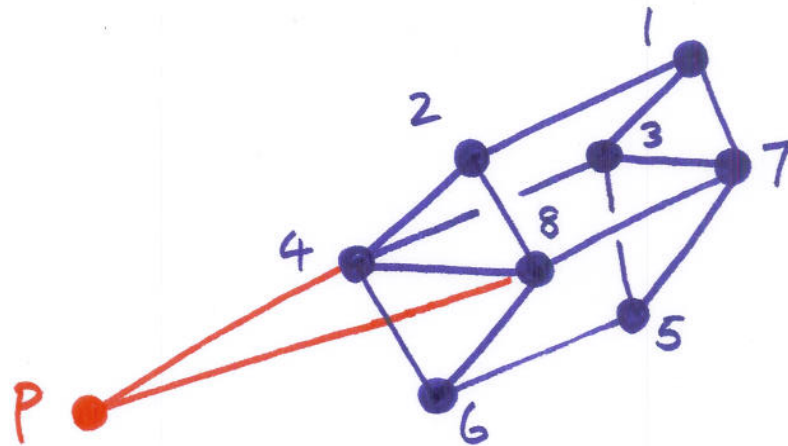
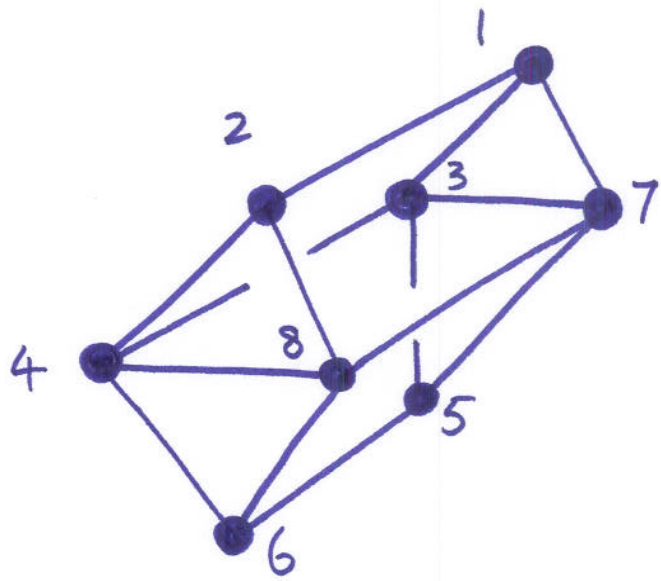
Resulting structures are polymatroids.

Example In a vector space V ,
let E be some set of subspaces
– if E consists of one-dim'l subspaces only,
then we have a matroid.

V_8 is not
representable

Proof

Suppose V_8 is
representable. Then it sits in a projective
space



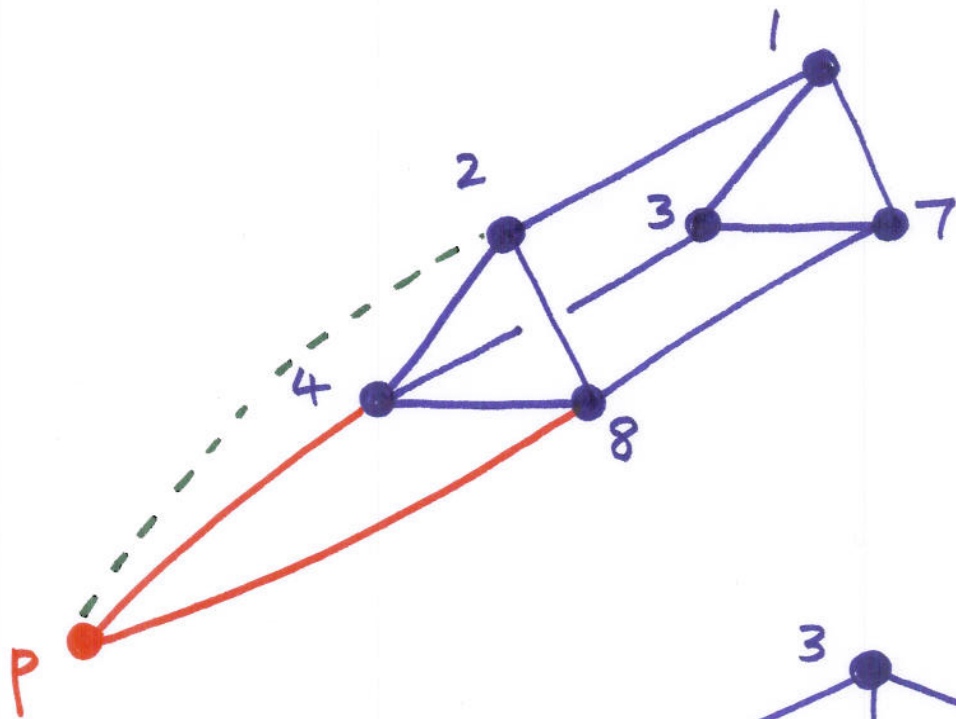
$$r(\{1, 2, 3, 4, p\}) + r(\{1, 2, 7, 8, p\})$$

$$\geq r(\{1, 2, 3, 4, 7, 8, p\}) + r(\{1, 2, p\})$$

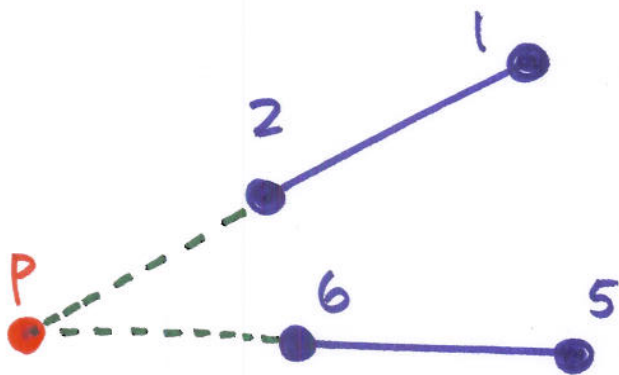
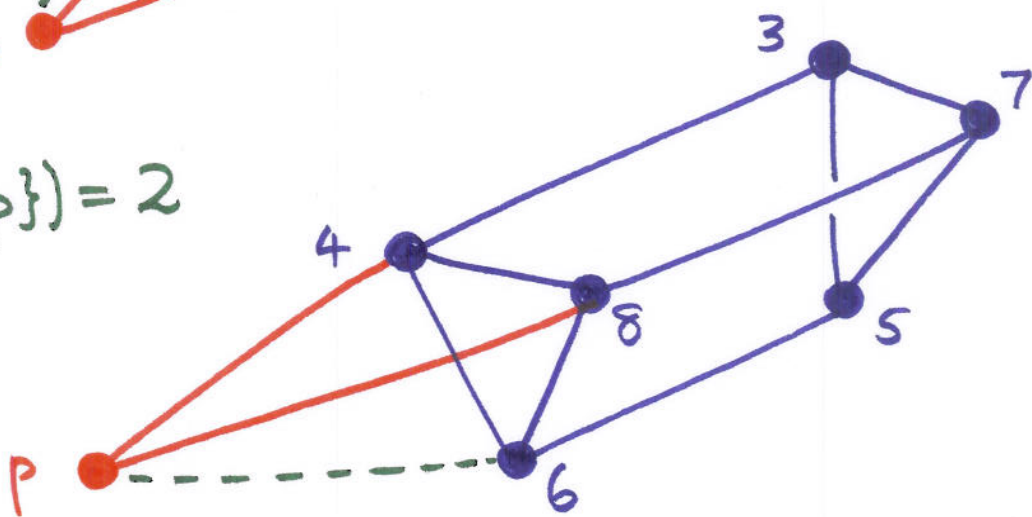
$$3 + 3 \geq 4 + r(\{1, 2, p\})$$

$$3 + 3 \geq 4 + r(\{1, 2, p\})$$

$$\text{so } r(\{1, 2, p\}) = 2.$$



$$r(\{5, 6, p\}) = 2$$



So $\{1, 2, 5, 6\}$ are
coplanar; not so in
 V_8 .



Summary

Basic matroid operations

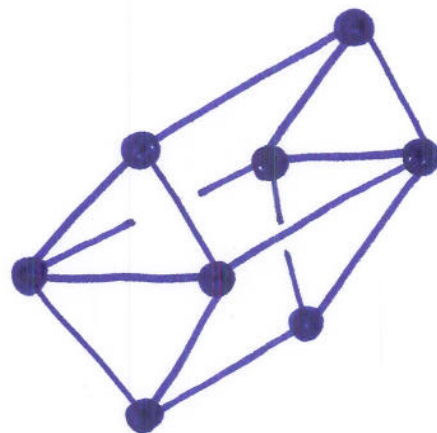
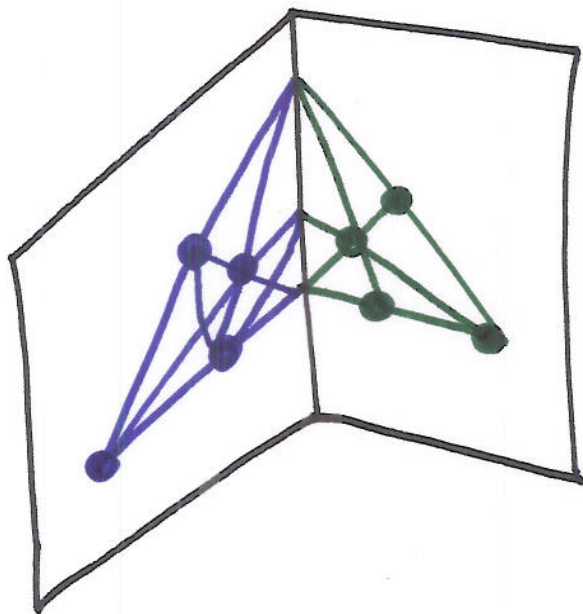
- duality
 - generalizes vector space orthogonality
 - matches up with coding theory duality
- deletion
- contraction
 - dual of deletion
 - corresponds geometrically to projection.
- minors
 - sequences of deletions and contractions

Basic example F -rep matroids

Rota's Conj. For all q , the set of excluded minors for $GF(q)$ -rep is finite

Non-representable matroids

Examples



V_8 , Vámos matroid

Conjecture (Mayhew, Newman, Welsh, Whittle)

Asymptotically almost every n -element matroid:

- is not representable
- is sparse paving
- has rank in $\{ \lfloor \frac{n}{2} \rfloor, \lceil \frac{n}{2} \rceil \}$
- has V_8 as a minor
- has any fixed sparse non-paving matroid as a minor